

人工智慧

(Artificial Intelligence)

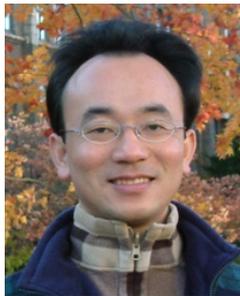
知識推理和知識表達

(Knowledge, Reasoning and Knowledge Representation)

1092AI04

MBA, IM, NTPU (M5010) (Spring 2021)

Wed 2, 3, 4 (9:10-12:00) (B8F40)



Min-Yuh Day

戴敏育

Associate Professor

副教授

Institute of Information Management, National Taipei University

國立臺北大學 資訊管理研究所

<https://web.ntpu.edu.tw/~myday>

2021-03-17



課程大綱 (Syllabus)

- | 週次 (Week) | 日期 (Date) | 內容 (Subject/Topics) |
|-----------|------------|--|
| 1 | 2021/02/24 | 人工智慧概論
(Introduction to Artificial Intelligence) |
| 2 | 2021/03/03 | 人工智慧和智慧代理人
(Artificial Intelligence and Intelligent Agents) |
| 3 | 2021/03/10 | 問題解決
(Problem Solving) |
| 4 | 2021/03/17 | 知識推理和知識表達
(Knowledge, Reasoning and Knowledge Representation) |
| 5 | 2021/03/24 | 不確定知識和推理
(Uncertain Knowledge and Reasoning) |
| 6 | 2021/03/31 | 人工智慧個案研究 I
(Case Study on Artificial Intelligence I) |

課程大綱 (Syllabus)

週次 (Week)	日期 (Date)	內容 (Subject/Topics)
7	2021/04/07	放假一天 (Day off)
8	2021/04/14	機器學習與監督式學習 (Machine Learning and Supervised Learning)
9	2021/04/21	期中報告 (Midterm Project Report)
10	2021/04/28	學習理論與綜合學習 (The Theory of Learning and Ensemble Learning)
11	2021/05/05	深度學習 (Deep Learning)
12	2021/05/12	人工智慧個案研究 II (Case Study on Artificial Intelligence II)

課程大綱 (Syllabus)

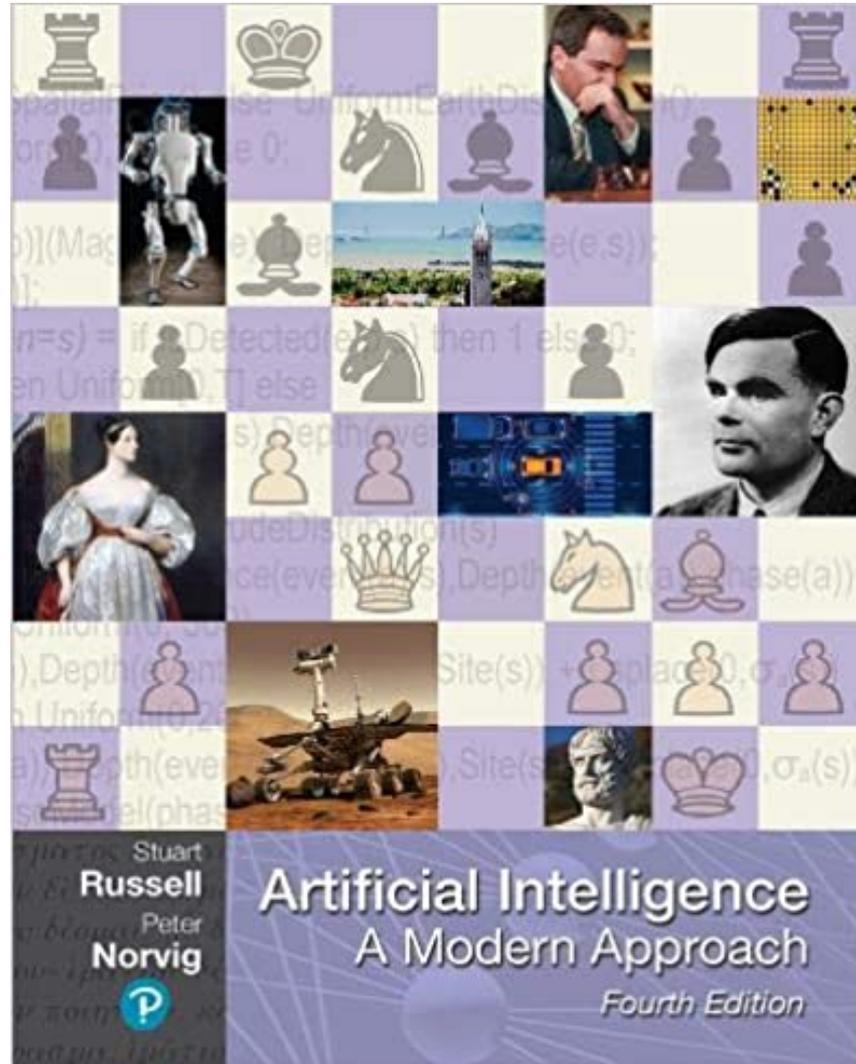
週次 (Week)	日期 (Date)	內容 (Subject/Topics)
13	2021/05/19	強化學習 (Reinforcement Learning)
14	2021/05/26	深度學習自然語言處理 (Deep Learning for Natural Language Processing)
15	2021/06/02	機器人技術 (Robotics)
16	2021/06/09	人工智慧哲學與倫理，人工智慧的未來 (Philosophy and Ethics of AI, The Future of AI)
17	2021/06/16	期末報告 I (Final Project Report I)
18	2021/06/23	期末報告 II (Final Project Report II)

**Knowledge,
Reasoning
and
Knowledge
Representation**

Outline

- **Logical Agents**
- **First-Order Logic**
- **Inference in First-Order Logic**
- **Knowledge Representation**
- **Automated Planning**

Stuart Russell and Peter Norvig (2020),
Artificial Intelligence: A Modern Approach,
4th Edition, Pearson



Source: Stuart Russell and Peter Norvig (2020), Artificial Intelligence: A Modern Approach, 4th Edition, Pearson

<https://www.amazon.com/Artificial-Intelligence-A-Modern-Approach/dp/0134610997/>

Artificial Intelligence: A Modern Approach

1. Artificial Intelligence
2. Problem Solving
3. Knowledge and Reasoning
4. Uncertain Knowledge and Reasoning
5. Machine Learning
6. Communicating, Perceiving, and Acting
7. Philosophy and Ethics of AI

Artificial Intelligence: Knowledge and Reasoning

Artificial Intelligence:

3. Knowledge and Reasoning

- Logical Agents
- First-Order Logic
- Inference in First-Order Logic
- Knowledge Representation
- Automated Planning

Intelligent Agents

4 Approaches of AI

2.

**Thinking Humanly:
The Cognitive
Modeling Approach**

3.

**Thinking Rationally:
The “Laws of Thought”
Approach**

1.

**Acting Humanly:
The Turing Test
Approach** (1950)

4.

**Acting Rationally:
The Rational Agent
Approach**

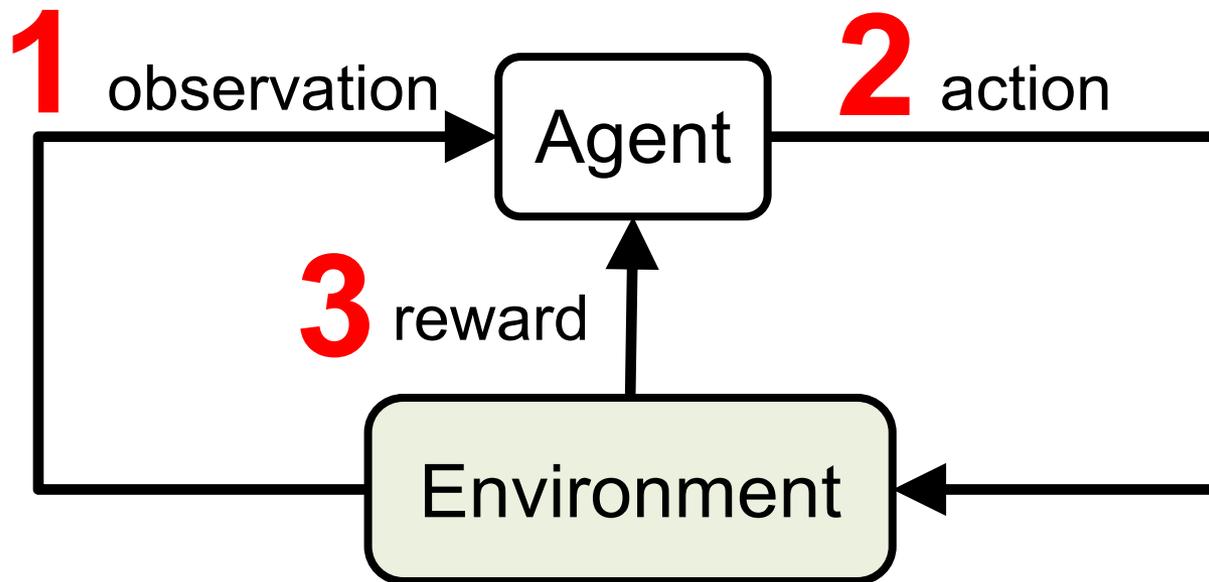
Reinforcement Learning (DL)

The diagram illustrates the Reinforcement Learning loop. It consists of two main components: an Agent and an Environment. The Agent is represented by a white rounded rectangle with a black border, positioned at the top. The Environment is represented by a light green rounded rectangle with a black border, positioned at the bottom. The Agent and Environment are connected by a vertical line, indicating the interaction between them.

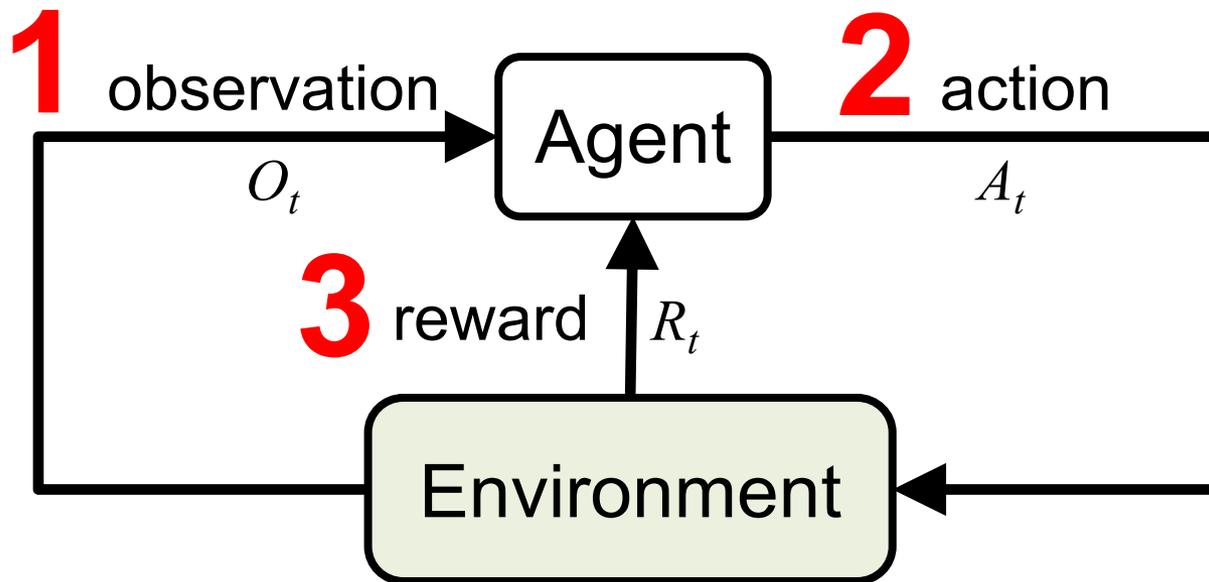
Agent

Environment

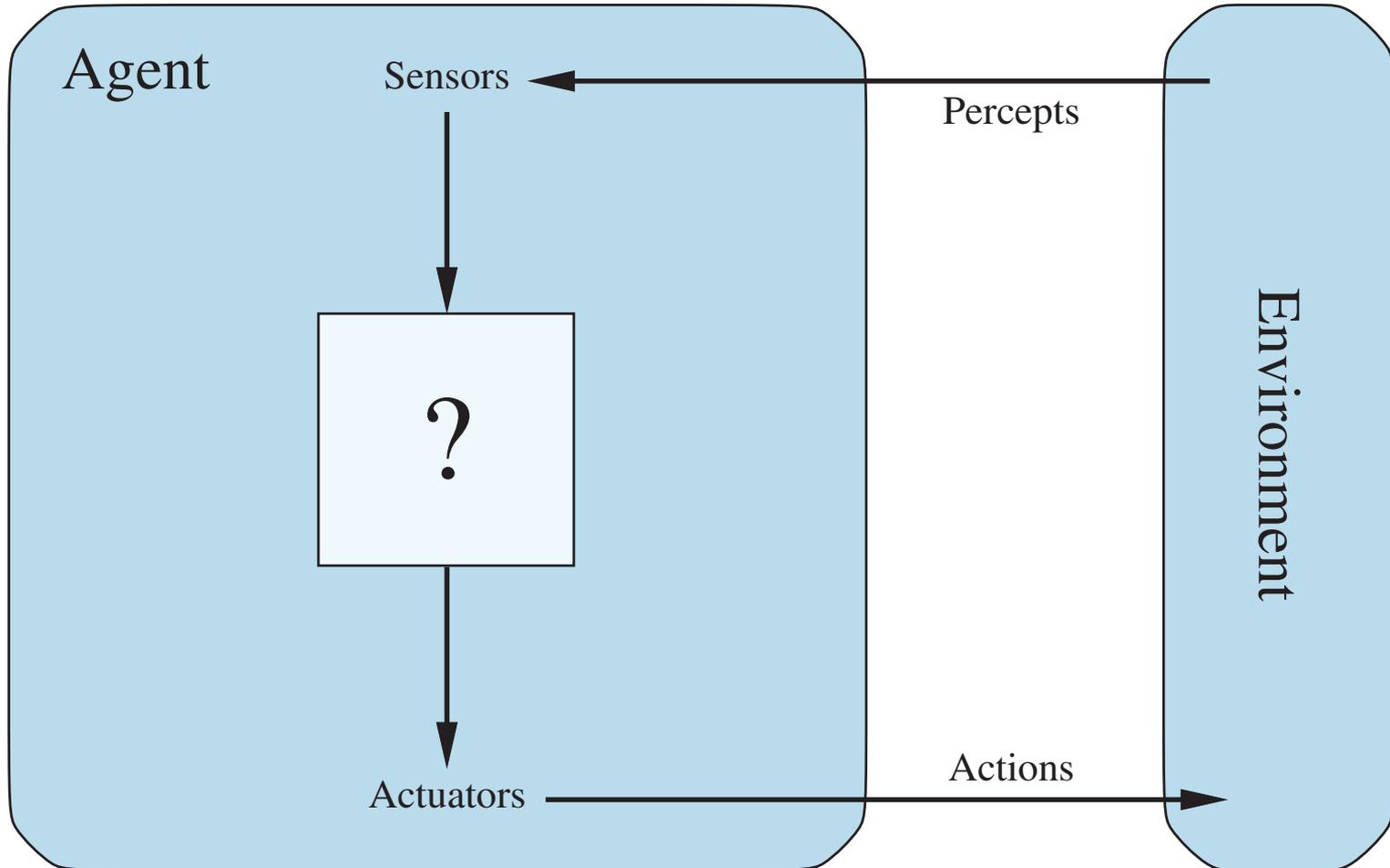
Reinforcement Learning (DL)



Reinforcement Learning (DL)



Agents interact with environments through sensors and actuators



Logical Agents

Logical Agents

Knowledge-based Agents

KB Agents

Knowledge-based Agent (KB Agent)

function KB-AGENT(*percept*) **returns** an *action*

persistent: *KB*, a knowledge base

t, a counter, initially 0, indicating time

TELL(*KB*, MAKE-PERCEPT-SENTENCE(*percept*, *t*))

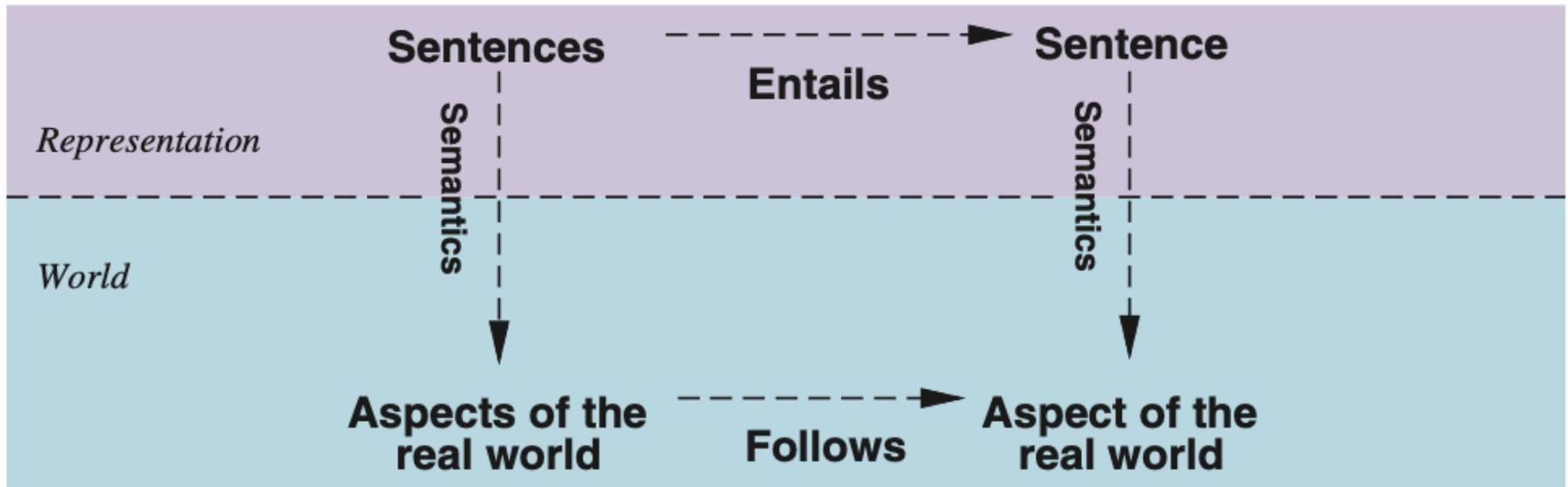
action \leftarrow ASK(*KB*, MAKE-ACTION-QUERY(*t*))

TELL(*KB*, MAKE-ACTION-SENTENCE(*action*, *t*))

t \leftarrow *t* + 1

return *action*

Sentences are physical configurations of the agent



Reasoning is a process of constructing new physical configurations from old ones

Logical reasoning should ensure that the new configurations represent aspects of the world that actually follow from the aspects that the old configurations represent.

A BNF (Backus–Naur Form) grammar of sentences in propositional logic

Sentence \rightarrow *AtomicSentence* | *ComplexSentence*

AtomicSentence \rightarrow *True* | *False* | *P* | *Q* | *R* | ...

ComplexSentence \rightarrow (*Sentence*)

| \neg *Sentence*

| *Sentence* \wedge *Sentence*

| *Sentence* \vee *Sentence*

| *Sentence* \Rightarrow *Sentence*

| *Sentence* \Leftrightarrow *Sentence*

OPERATOR PRECEDENCE : $\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$

Truth Tables (TT) for the Five Logical Connectives

P	Q	$\neg P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
<i>false</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>
<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>
<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>
<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>

A Truth Table constructed for the knowledge base given in the text

$B_{1,1}$	$B_{2,1}$	$P_{1,1}$	$P_{1,2}$	$P_{2,1}$	$P_{2,2}$	$P_{3,1}$	R_1	R_2	R_3	R_4	R_5	KB
<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	<i>false</i>						
<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>
\vdots												
<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>
<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<u><i>true</i></u>
<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<u><i>true</i></u>
<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<u><i>true</i></u>						
<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>
\vdots												
<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>						

A Truth-Table (TT) enumeration algorithm for deciding propositional entailment

function TT-ENTAILS?(KB, α) **returns** *true* or *false*

inputs: KB , the knowledge base, a sentence in propositional logic
 α , the query, a sentence in propositional logic

symbols \leftarrow a list of the proposition symbols in KB and α

return TT-CHECK-ALL($KB, \alpha, symbols, \{ \}$)

function TT-CHECK-ALL($KB, \alpha, symbols, model$) **returns** *true* or *false*

if EMPTY?(*symbols*) **then**

if PL-TRUE?($KB, model$) **then return** PL-TRUE?($\alpha, model$)

else return *true* // when KB is false, always return *true*

else

$P \leftarrow$ FIRST(*symbols*)

rest \leftarrow REST(*symbols*)

return (TT-CHECK-ALL($KB, \alpha, rest, model \cup \{P = true\}$)

and

 TT-CHECK-ALL($KB, \alpha, rest, model \cup \{P = false\}$))

Standard Logical Equivalences

The symbols α , β , and γ stand for arbitrary sentences of propositional logic.

$$(\alpha \wedge \beta) \equiv (\beta \wedge \alpha) \quad \text{commutativity of } \wedge$$

$$(\alpha \vee \beta) \equiv (\beta \vee \alpha) \quad \text{commutativity of } \vee$$

$$((\alpha \wedge \beta) \wedge \gamma) \equiv (\alpha \wedge (\beta \wedge \gamma)) \quad \text{associativity of } \wedge$$

$$((\alpha \vee \beta) \vee \gamma) \equiv (\alpha \vee (\beta \vee \gamma)) \quad \text{associativity of } \vee$$

$$\neg(\neg\alpha) \equiv \alpha \quad \text{double-negation elimination}$$

$$(\alpha \Rightarrow \beta) \equiv (\neg\beta \Rightarrow \neg\alpha) \quad \text{contraposition}$$

$$(\alpha \Rightarrow \beta) \equiv (\neg\alpha \vee \beta) \quad \text{implication elimination}$$

$$(\alpha \Leftrightarrow \beta) \equiv ((\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)) \quad \text{biconditional elimination}$$

$$\neg(\alpha \wedge \beta) \equiv (\neg\alpha \vee \neg\beta) \quad \text{De Morgan}$$

$$\neg(\alpha \vee \beta) \equiv (\neg\alpha \wedge \neg\beta) \quad \text{De Morgan}$$

$$(\alpha \wedge (\beta \vee \gamma)) \equiv ((\alpha \wedge \beta) \vee (\alpha \wedge \gamma)) \quad \text{distributivity of } \wedge \text{ over } \vee$$

$$(\alpha \vee (\beta \wedge \gamma)) \equiv ((\alpha \vee \beta) \wedge (\alpha \vee \gamma)) \quad \text{distributivity of } \vee \text{ over } \wedge$$

A grammar for Conjunctive Normal Form (CNF), Horn clauses, and definite clauses

CNFSentence \rightarrow *Clause*₁ $\wedge \dots \wedge$ *Clause*_n

Clause \rightarrow *Literal*₁ $\vee \dots \vee$ *Literal*_m

Fact \rightarrow *Symbol*

Literal \rightarrow *Symbol* | \neg *Symbol*

Symbol \rightarrow *P* | *Q* | *R* | ...

HornClauseForm \rightarrow *DefiniteClauseForm* | *GoalClauseForm*

DefiniteClauseForm \rightarrow *Fact* | (*Symbol*₁ $\wedge \dots \wedge$ *Symbol*_l) \Rightarrow *Symbol*

GoalClauseForm \rightarrow (*Symbol*₁ $\wedge \dots \wedge$ *Symbol*_l) \Rightarrow *False*

A simple resolution algorithm for propositional logic

function PL-RESOLUTION(KB, α) **returns** *true* or *false*
inputs: KB , the knowledge base, a sentence in propositional logic
 α , the query, a sentence in propositional logic

$clauses \leftarrow$ the set of clauses in the CNF representation of $KB \wedge \neg\alpha$
 $new \leftarrow \{ \}$

while *true* **do**
 for each pair of clauses C_i, C_j **in** $clauses$ **do**
 $resolvents \leftarrow$ PL-RESOLVE(C_i, C_j)
 if $resolvents$ contains the empty clause **then return** *true*
 $new \leftarrow new \cup resolvents$

if $new \subseteq clauses$ **then return** *false*
 $clauses \leftarrow clauses \cup new$

The forward-chaining algorithm for propositional logic

function PL-FC-ENTAILS?(KB, q) **returns** *true* or *false*

inputs: KB , the knowledge base, a set of propositional definite clauses

q , the query, a proposition symbol

$count \leftarrow$ a table, where $count[c]$ is initially the number of symbols in clause c 's premise

$inferred \leftarrow$ a table, where $inferred[s]$ is initially *false* for all symbols

$queue \leftarrow$ a queue of symbols, initially symbols known to be true in KB

while $queue$ is not empty **do**

$p \leftarrow \text{POP}(queue)$

if $p = q$ **then return** *true*

if $inferred[p] = false$ **then**

$inferred[p] \leftarrow true$

for each clause c in KB where p is in c .PREMISE **do**

decrement $count[c]$

if $count[c] = 0$ **then** add c .CONCLUSION to $queue$

return *false*

A set of Horn clauses

$$P \Rightarrow Q$$

$$L \wedge M \Rightarrow P$$

$$B \wedge L \Rightarrow M$$

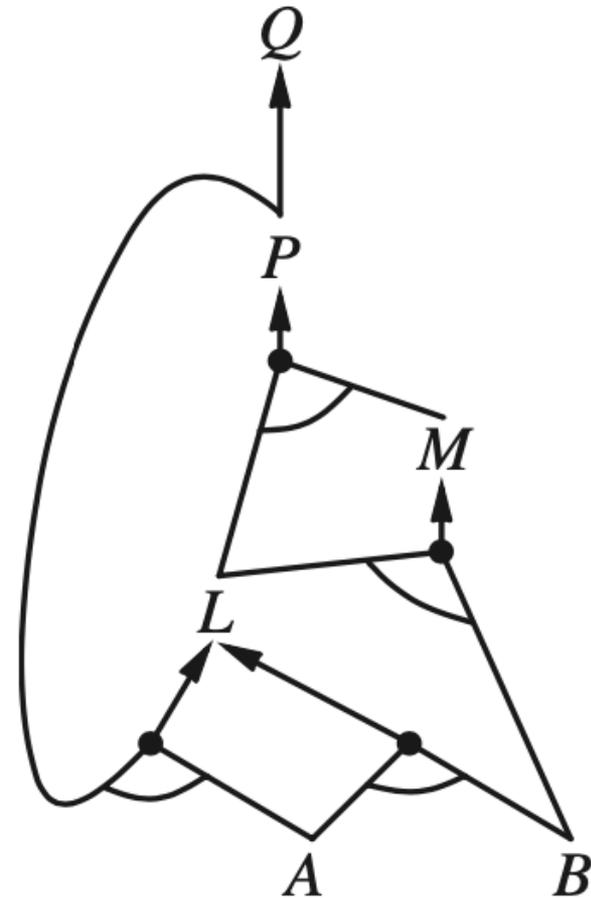
$$A \wedge P \Rightarrow L$$

$$A \wedge B \Rightarrow L$$

A

B

(a)



(b)

The corresponding AND–OR graph

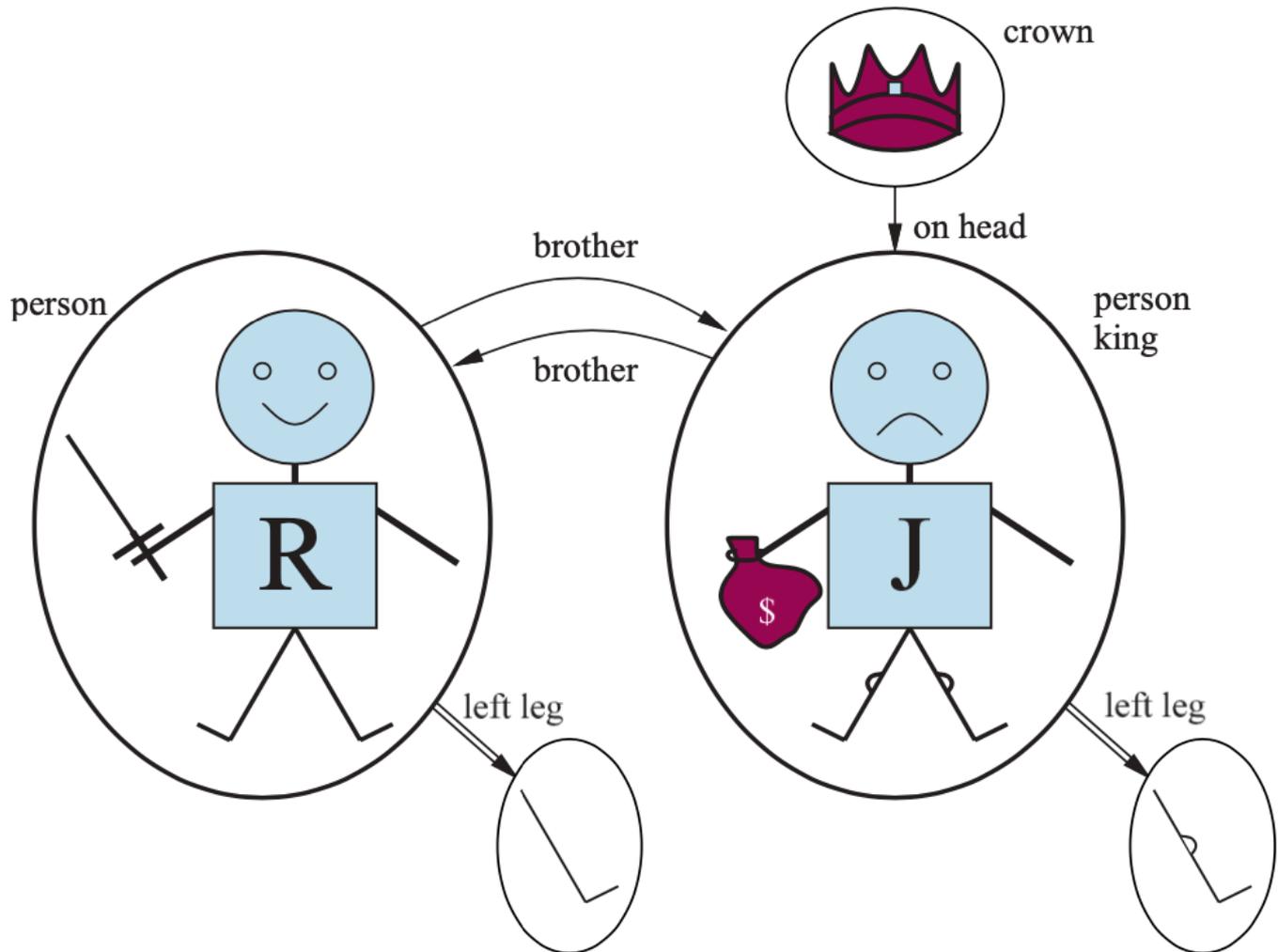
First-Order Logic

Formal languages and their ontological and epistemological commitments

Language	Ontological Commitment (What exists in the world)	Epistemological Commitment (What an agent believes about facts)
Propositional logic	facts	true/false/unknown
First-order logic	facts, objects, relations	true/false/unknown
Temporal logic	facts, objects, relations, times	true/false/unknown
Probability theory	facts	degree of belief $\in [0, 1]$
Fuzzy logic	facts with degree of truth $\in [0, 1]$	known interval value

A model containing five objects

two binary relations (brother and on-head), three unary relations (person, king, and crown), and one unary function (left-leg).



The syntax of first-order logic with equality

Sentence → *AtomicSentence* | *ComplexSentence*

AtomicSentence → *Predicate* | *Predicate*(*Term*, ...) | *Term* = *Term*

ComplexSentence → (*Sentence*)

| \neg *Sentence*

| *Sentence* \wedge *Sentence*

| *Sentence* \vee *Sentence*

| *Sentence* \Rightarrow *Sentence*

| *Sentence* \Leftrightarrow *Sentence*

| *Quantifier* *Variable*, ... *Sentence*

Term → *Function*(*Term*, ...)

| *Constant*

| *Variable*

Quantifier → \forall | \exists

Constant → *A* | *X*₁ | *John* | ...

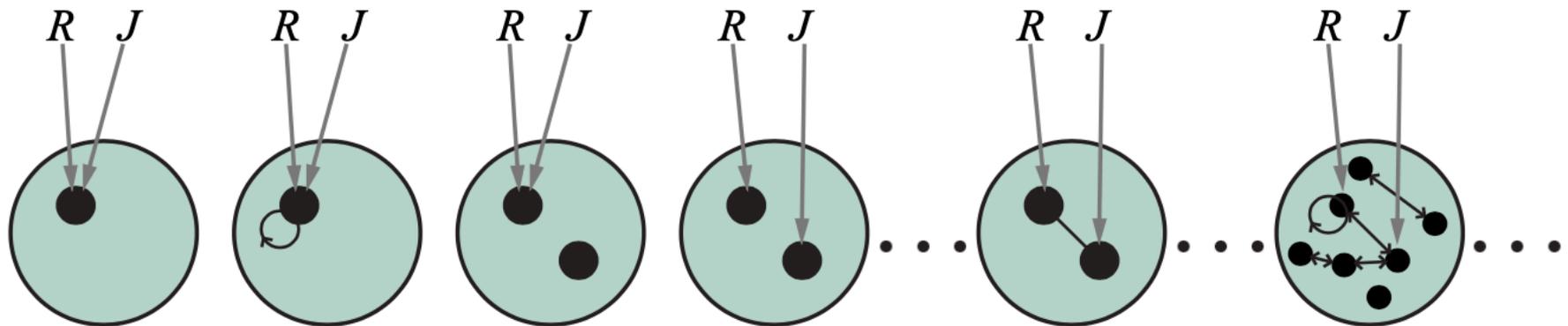
Variable → *a* | *x* | *s* | ...

Predicate → *True* | *False* | *After* | *Loves* | *Raining* | ...

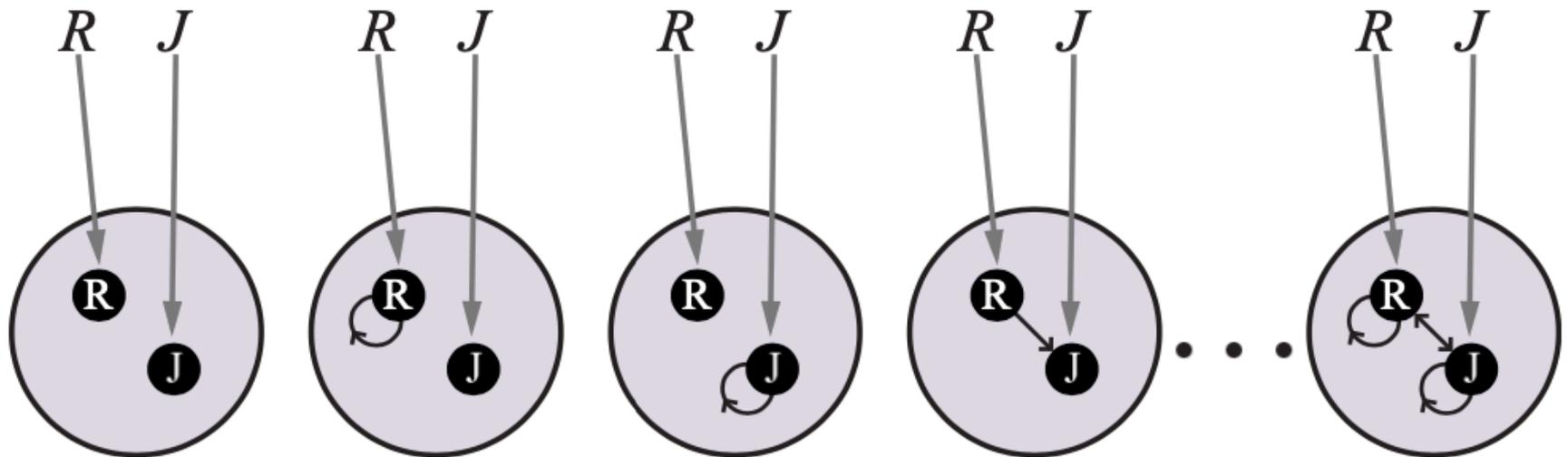
Function → *Mother* | *LeftLeg* | ...

OPERATOR PRECEDENCE : $\neg, =, \wedge, \vee, \Rightarrow, \Leftrightarrow$

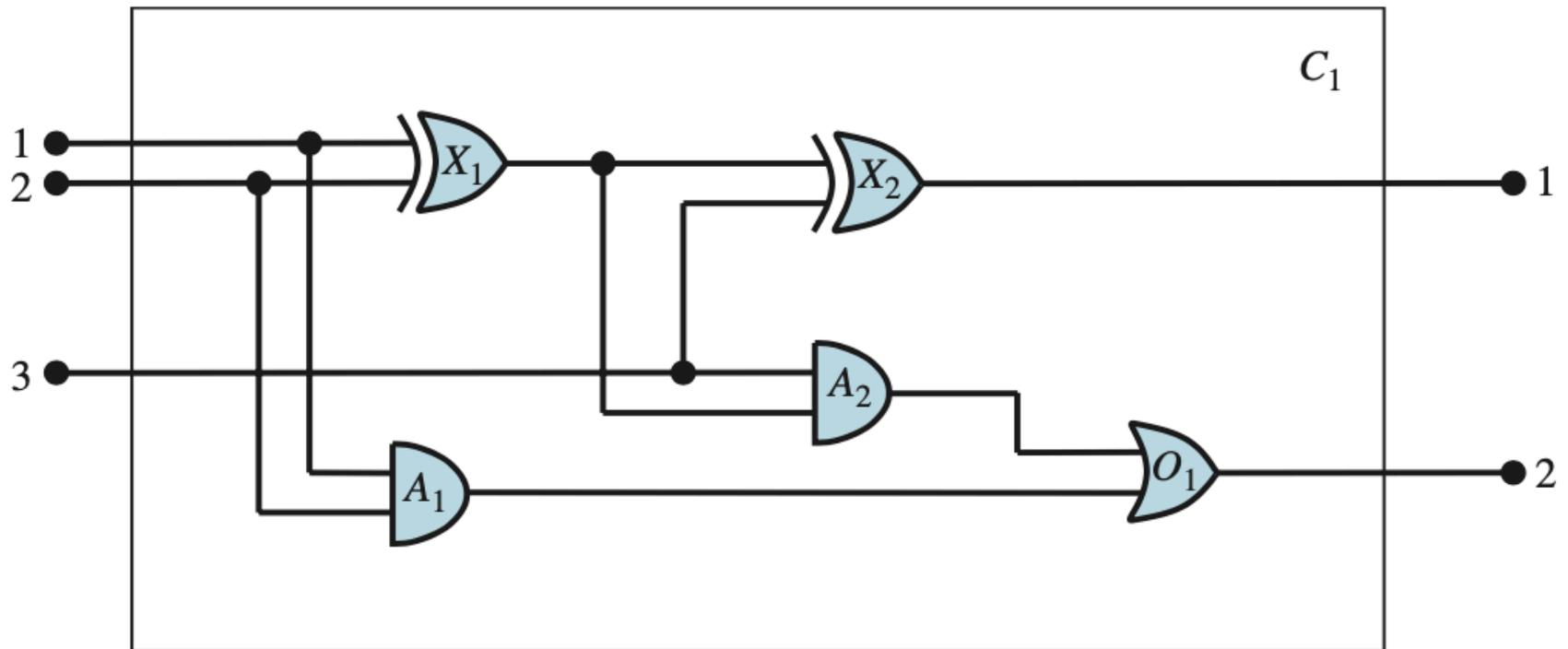
Some members of the set of all models for a language with two constant symbols, R and J , and one binary relation symbol



Some members of the set of all models for a language with two constant symbols, R and J , and one binary relation symbol, under database semantics



A digital circuit C₁, purporting to be a one-bit full adder.



Inference in First-Order Logic

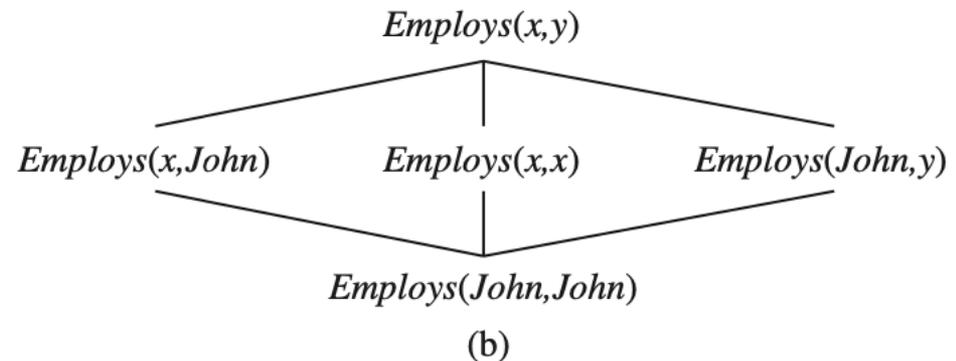
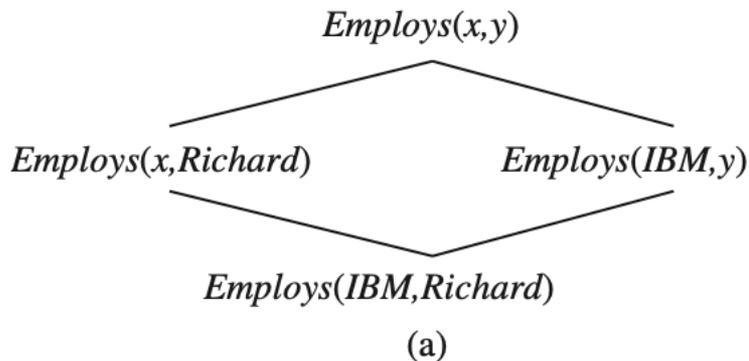
The unification algorithm

function UNIFY($x, y, \theta = \text{empty}$) **returns** a substitution to make x and y identical, or *failure*
if $\theta = \text{failure}$ **then return** *failure*
else if $x = y$ **then return** θ
else if VARIABLE?(x) **then return** UNIFY-VAR(x, y, θ)
else if VARIABLE?(y) **then return** UNIFY-VAR(y, x, θ)
else if COMPOUND?(x) **and** COMPOUND?(y) **then**
 return UNIFY(ARGS(x), ARGS(y), UNIFY(OP(x), OP(y), θ))
else if LIST?(x) **and** LIST?(y) **then**
 return UNIFY(REST(x), REST(y), UNIFY(FIRST(x), FIRST(y), θ))
else return *failure*

function UNIFY-VAR(var, x, θ) **returns** a substitution
if $\{var/val\} \in \theta$ for some val **then return** UNIFY(val, x, θ)
else if $\{x/val\} \in \theta$ for some val **then return** UNIFY(var, val, θ)
else if OCCUR-CHECK?(var, x) **then return** *failure*
else return add $\{var/x\}$ to θ

The subsumption lattice whose lowest node is *Employs (IBM , Richard)*

The subsumption lattice for the sentence *Employs (John, John)*



A conceptually straightforward, but inefficient, forward-chaining algorithm

function FOL-FC-ASK(KB, α) **returns** a substitution or *false*

inputs: KB , the knowledge base, a set of first-order definite clauses

α , the query, an atomic sentence

while *true* **do**

$new \leftarrow \{ \}$ // *The set of new sentences inferred on each iteration*

for each *rule* **in** KB **do**

$(p_1 \wedge \dots \wedge p_n \Rightarrow q) \leftarrow \text{STANDARDIZE-VARIABLES}(\text{rule})$

for each θ such that $\text{SUBST}(\theta, p_1 \wedge \dots \wedge p_n) = \text{SUBST}(\theta, p'_1 \wedge \dots \wedge p'_n)$

for some p'_1, \dots, p'_n in KB

$q' \leftarrow \text{SUBST}(\theta, q)$

if q' does not unify with some sentence already in KB or *new* **then**

add q' to *new*

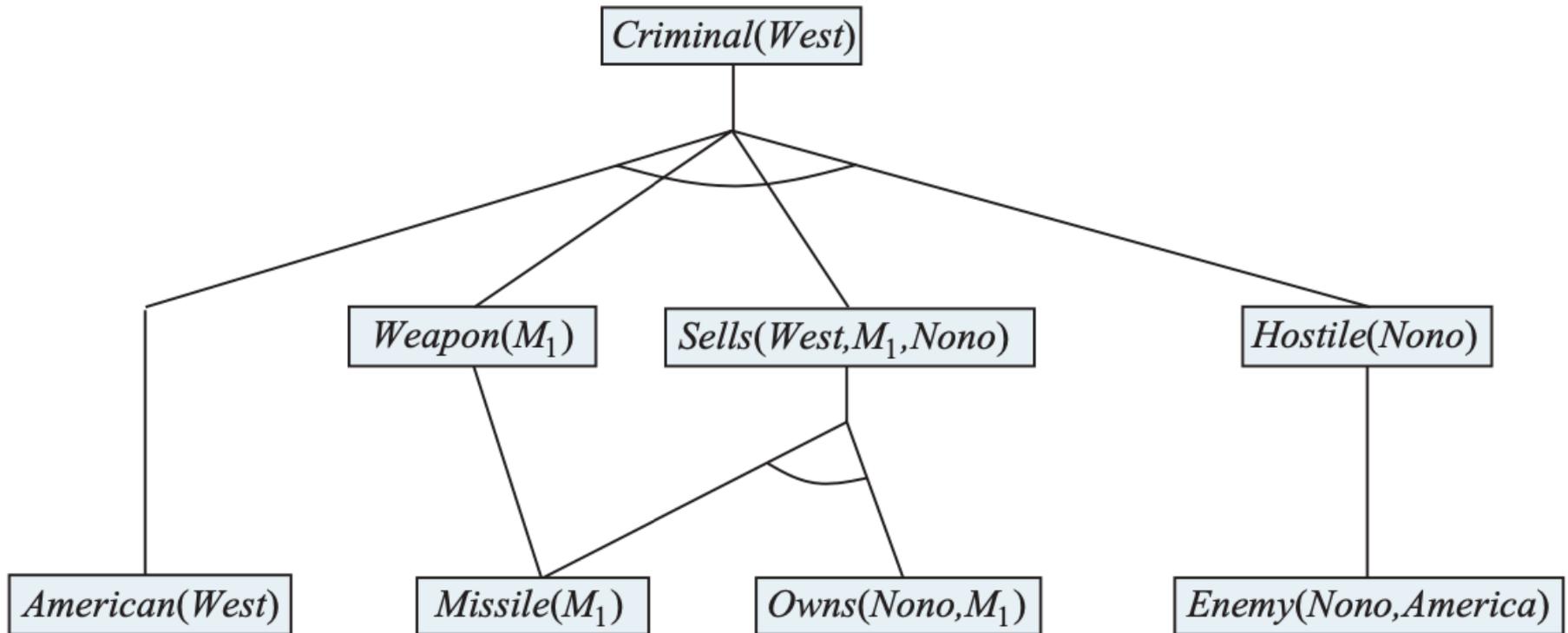
$\phi \leftarrow \text{UNIFY}(q', \alpha)$

if ϕ is not *failure* **then return** ϕ

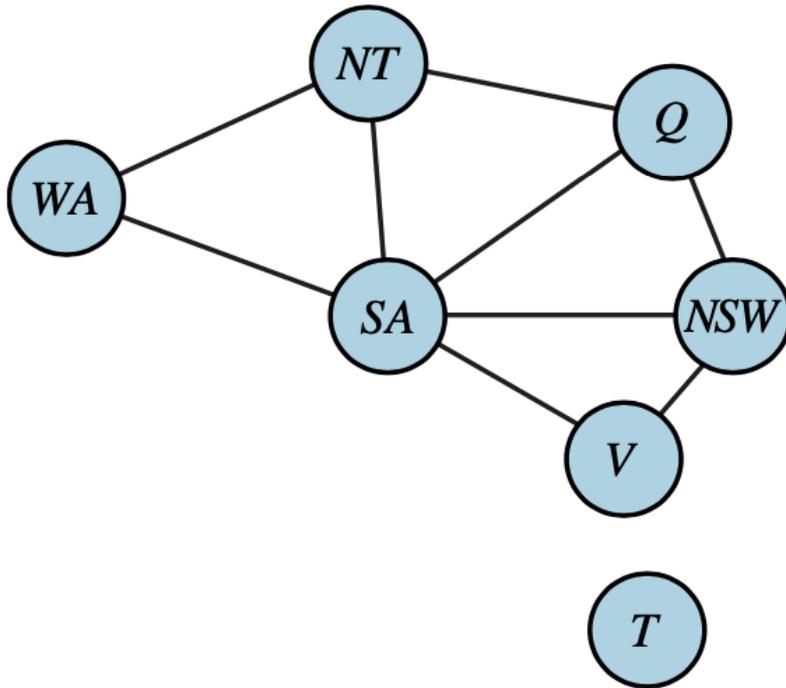
if $new = \{ \}$ **then return** *false*

add *new* to KB

The proof tree generated by forward chaining on the crime example



Constraint graph for coloring the map of Australia



(a)

$$\begin{aligned} & Diff(wa, nt) \wedge Diff(wa, sa) \wedge \\ & Diff(nt, q) \wedge Diff(nt, sa) \wedge \\ & Diff(q, nsw) \wedge Diff(q, sa) \wedge \\ & Diff(nsw, v) \wedge Diff(nsw, sa) \wedge \\ & Diff(v, sa) \Rightarrow Colorable() \end{aligned}$$

$$\begin{aligned} & Diff(Blue, Red) \quad Diff(Blue, Green) \\ & Diff(Green, Red) \quad Diff(Green, Blue) \\ & Diff(Red, Blue) \quad Diff(Red, Green) \end{aligned}$$

(b)

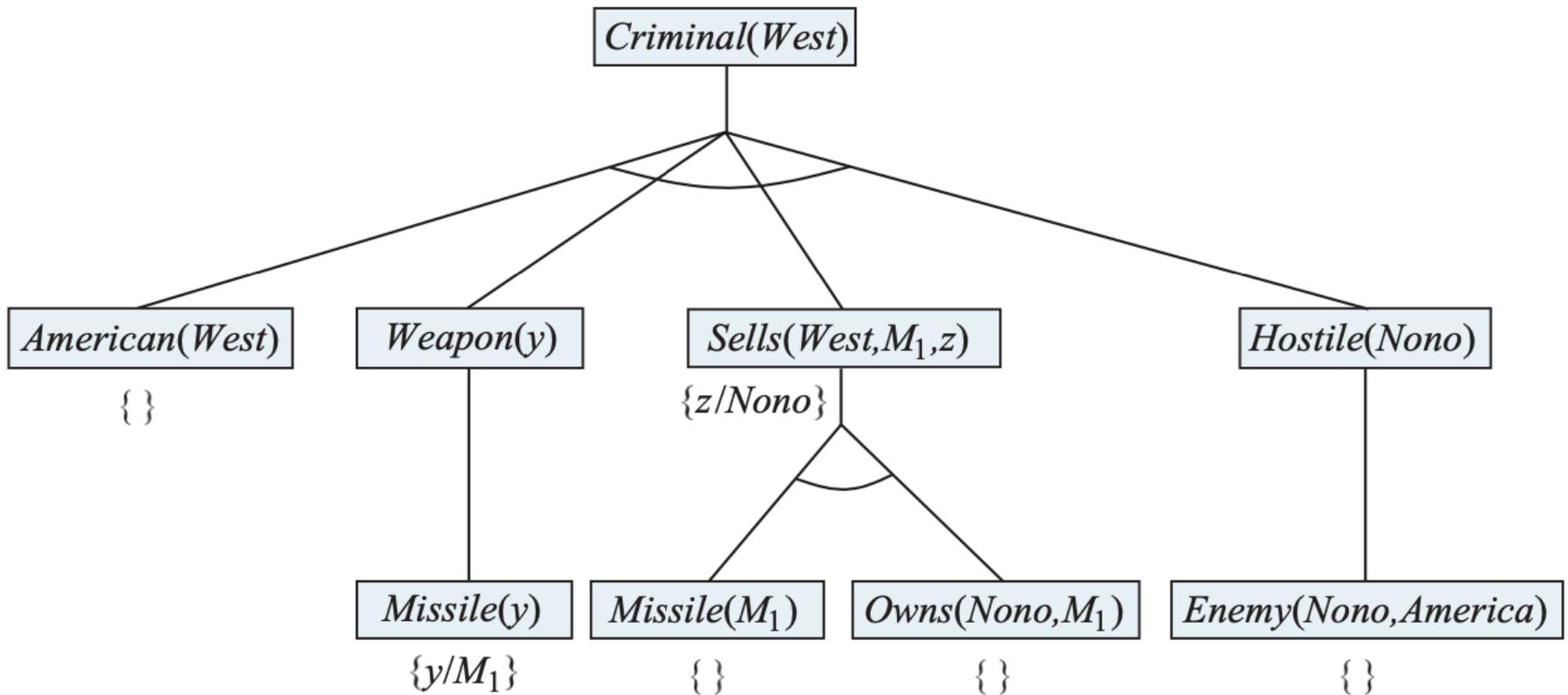
A simple backward-chaining algorithm for first-order knowledge bases

function FOL-BC-ASK($KB, query$) **returns** a generator of substitutions
return FOL-BC-OR($KB, query, \{ \}$)

function FOL-BC-OR($KB, goal, \theta$) **returns** a substitution
for each $rule$ in FETCH-RULES-FOR-GOAL($KB, goal$) **do**
 $(lhs \Rightarrow rhs) \leftarrow$ STANDARDIZE-VARIABLES($rule$)
 for each θ' in FOL-BC-AND($KB, lhs, UNIFY(rhs, goal, \theta)$) **do**
 yield θ'

function FOL-BC-AND($KB, goals, \theta$) **returns** a substitution
if $\theta = failure$ **then return**
else if LENGTH($goals$) = 0 **then yield** θ
else
 $first, rest \leftarrow$ FIRST($goals$), REST($goals$)
 for each θ' in FOL-BC-OR($KB, SUBST(\theta, first), \theta$) **do**
 for each θ'' in FOL-BC-AND($KB, rest, \theta'$) **do**
 yield θ''

Proof tree constructed by backward chaining to prove that West is a criminal



Pseudocode representing the result of compiling the Append predicate

procedure APPEND($ax, y, az, continuation$)

$trail \leftarrow$ GLOBAL-TRAIL-POINTER()

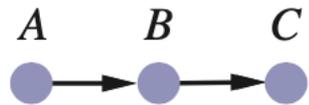
if $ax = []$ and UNIFY(y, az) **then** CALL($continuation$)

RESET-TRAIL($trail$)

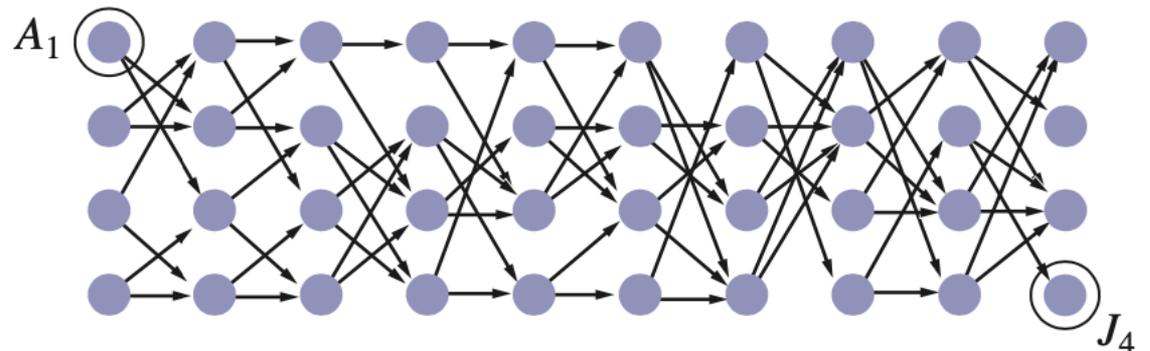
$a, x, z \leftarrow$ NEW-VARIABLE(), NEW-VARIABLE(), NEW-VARIABLE()

if UNIFY($ax, [a] + x$) and UNIFY($az, [a | z]$) **then** APPEND($x, y, z, continuation$)

Finding a path from A to C can lead Prolog into an infinite loop.

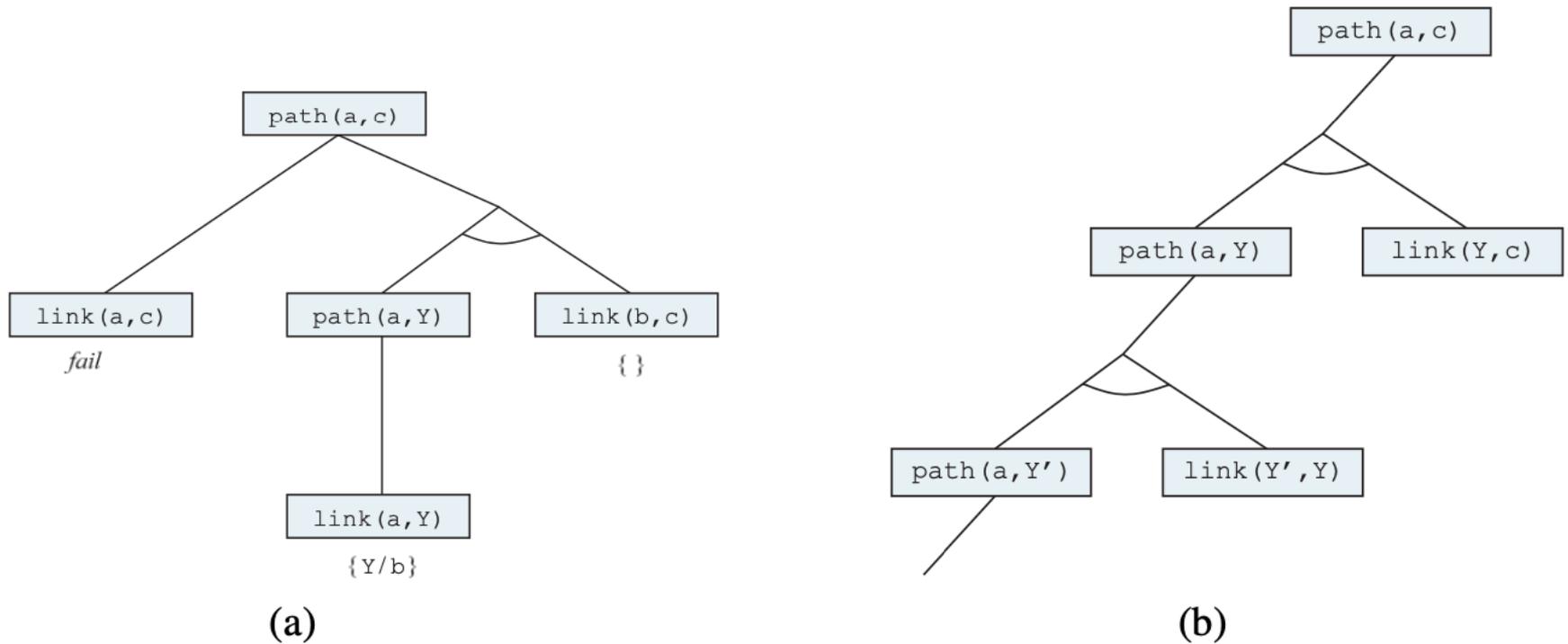


(a)



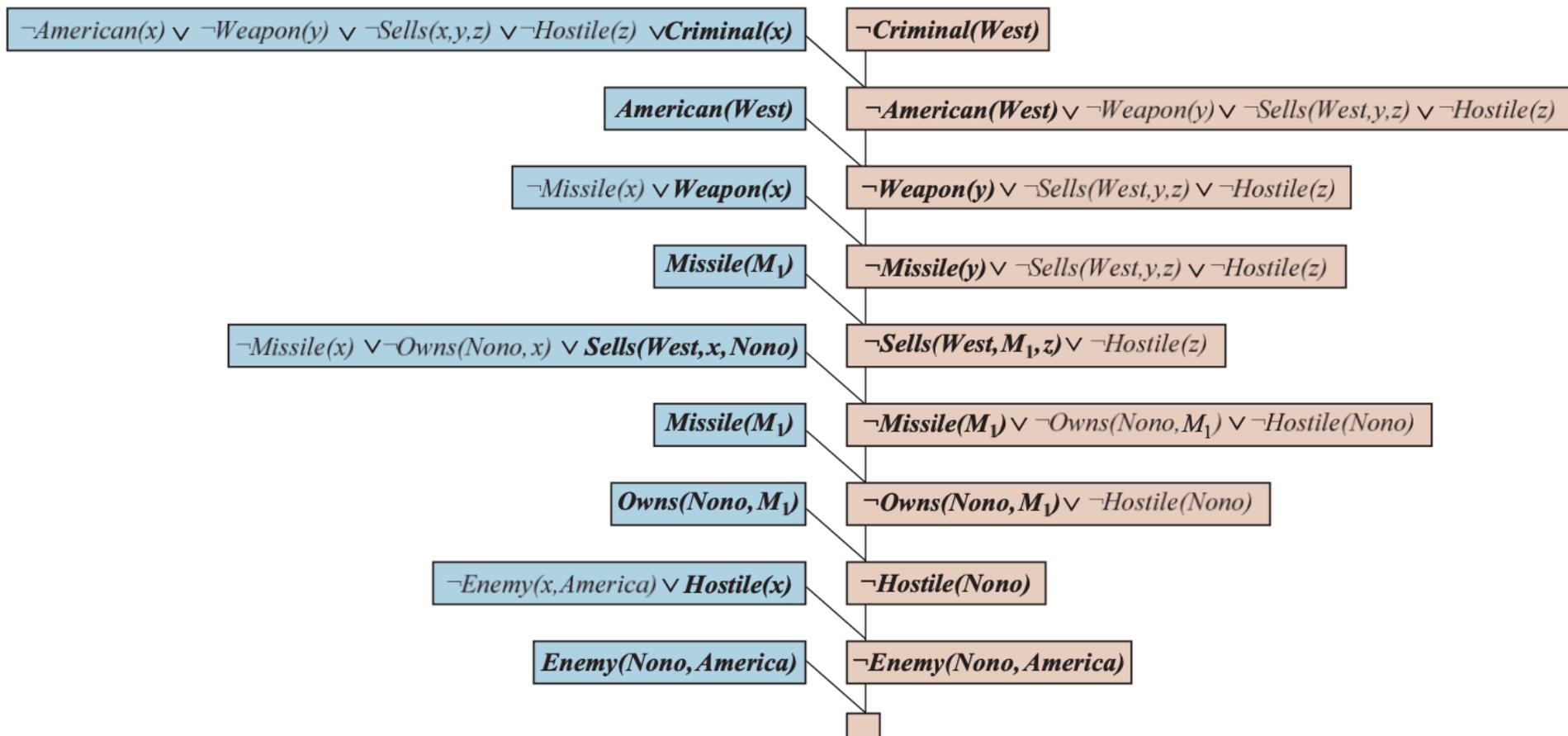
(b)

Proof that a path exists from A to C.

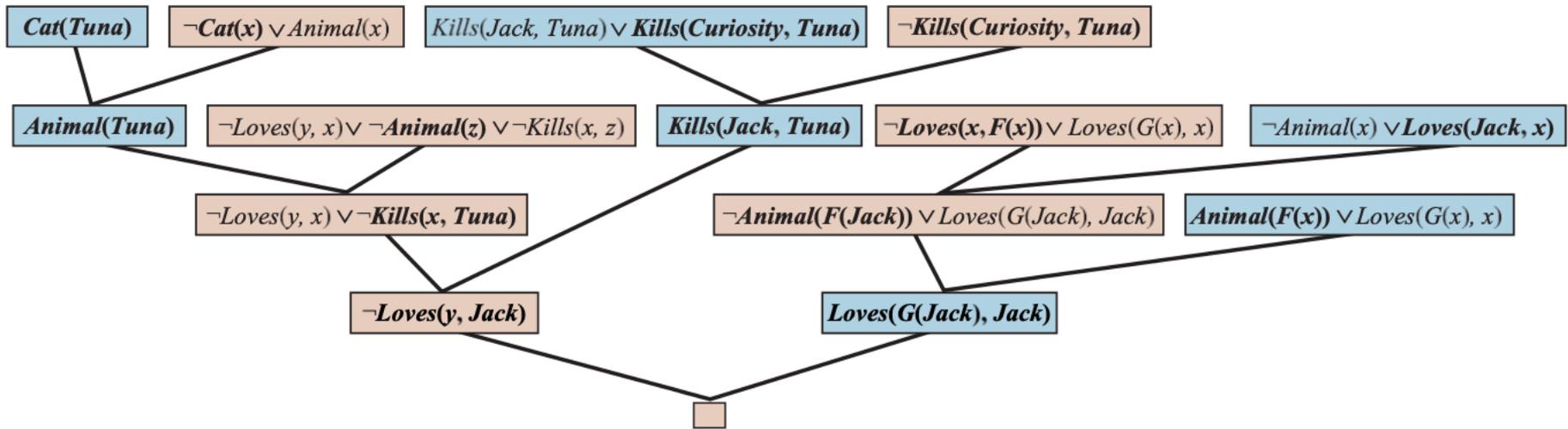


Infinite proof tree generated when the clauses are in the “wrong” order

A resolution proof that West is a criminal



A resolution proof that Curiosity killed the cat



Structure of a completeness proof for resolution

Any set of sentences S is representable in clausal form

Assume S is unsatisfiable, and in clausal form

Some set S' of ground instances is unsatisfiable

Resolution can find a contradiction in S'

There is a resolution proof for the contradiction in S'

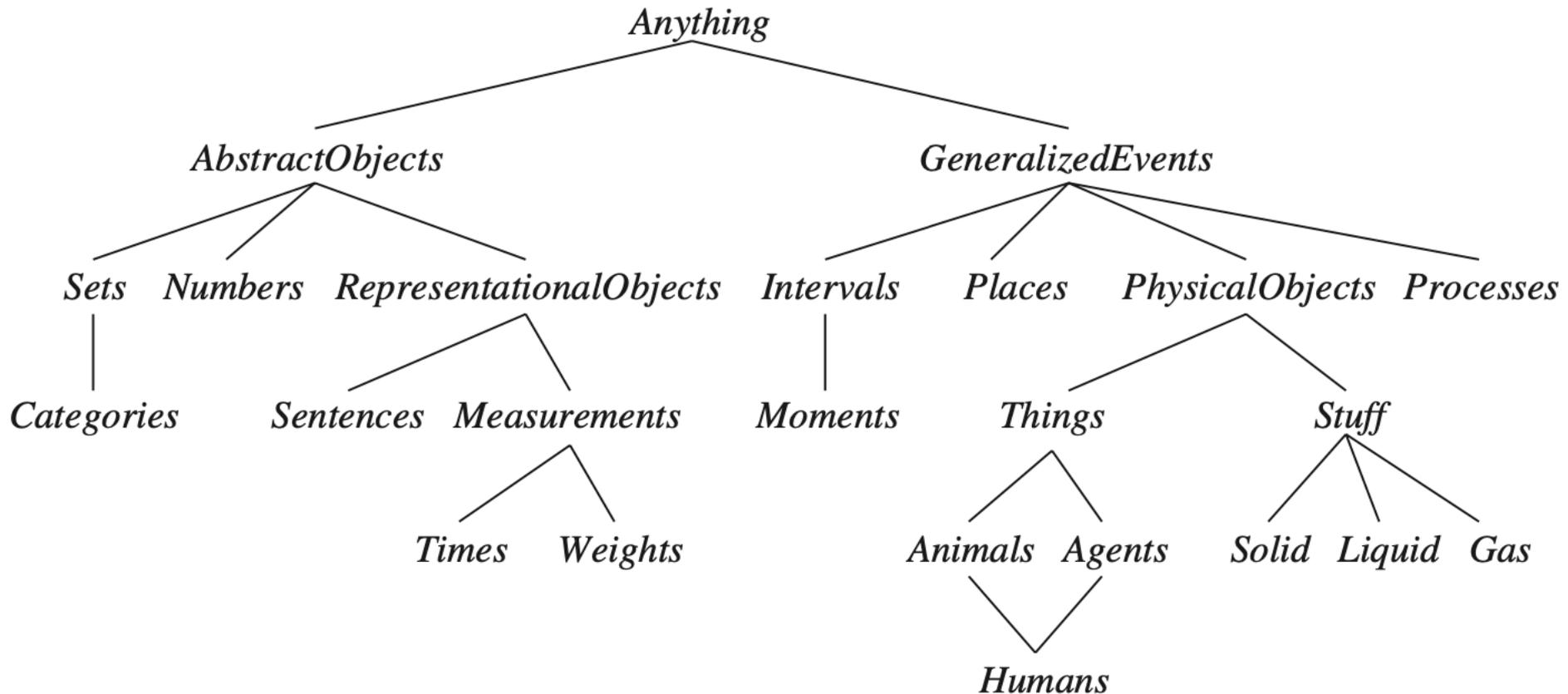
Herbrand's theorem

Ground resolution theorem

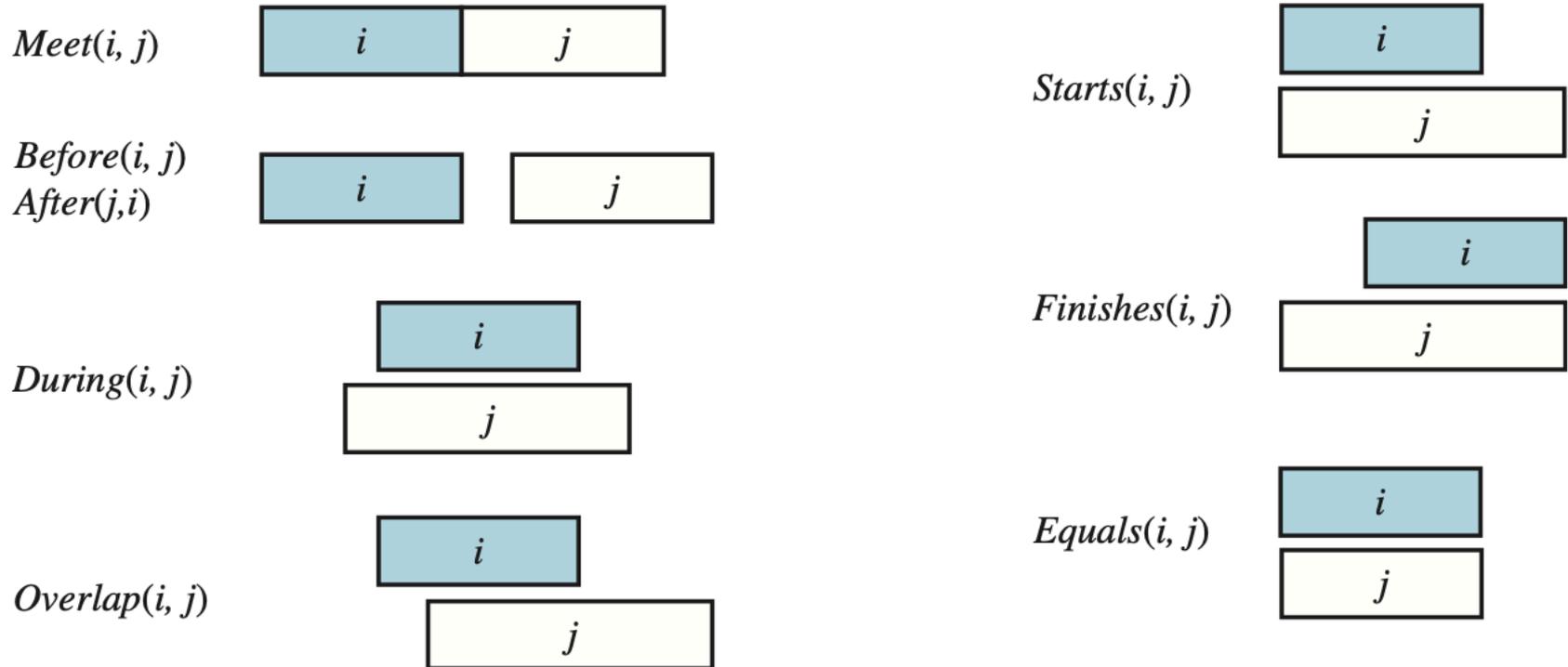
Lifting lemma

Knowledge Representation

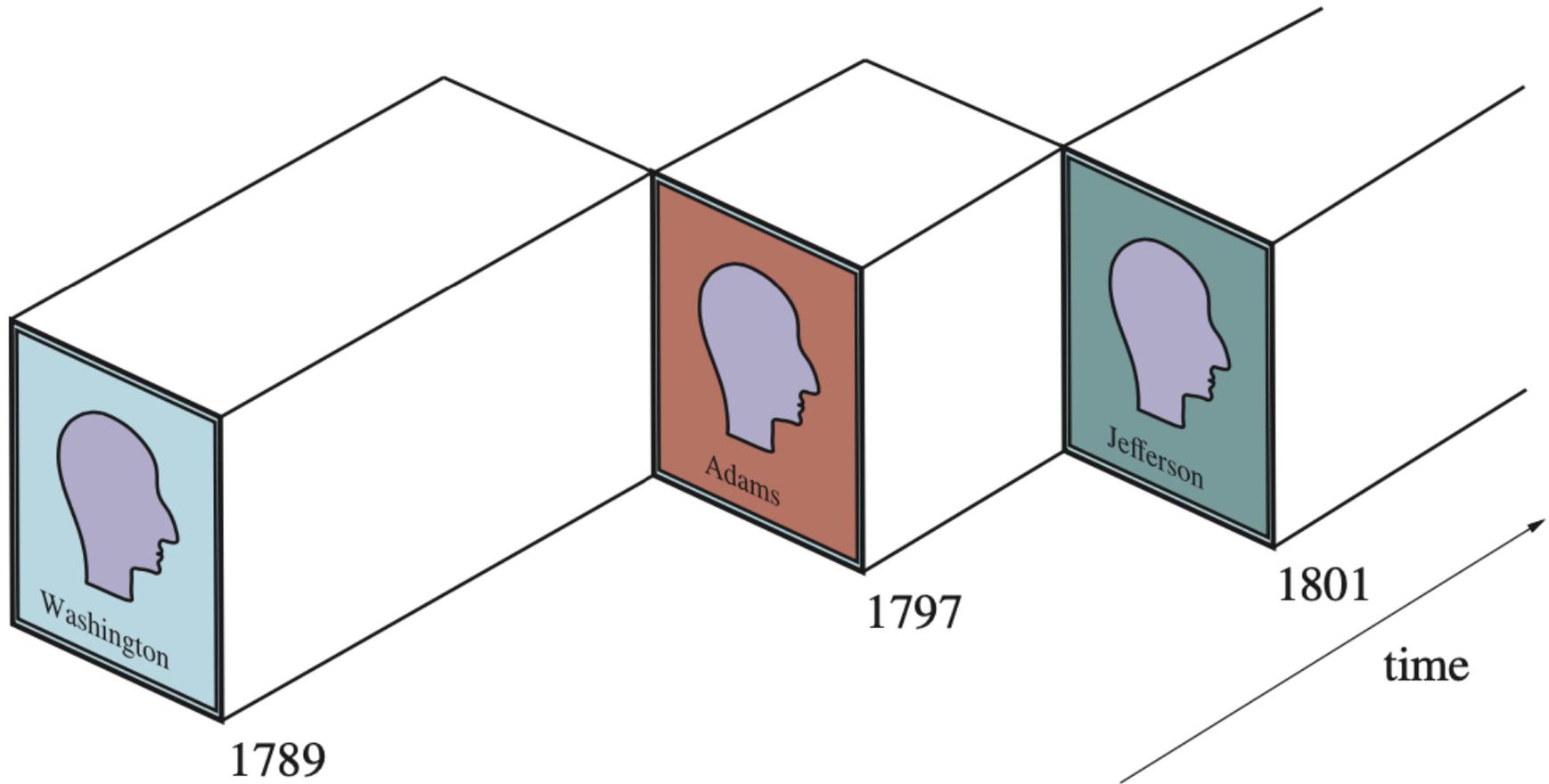
The Upper Ontology of the World



Predicates on time intervals

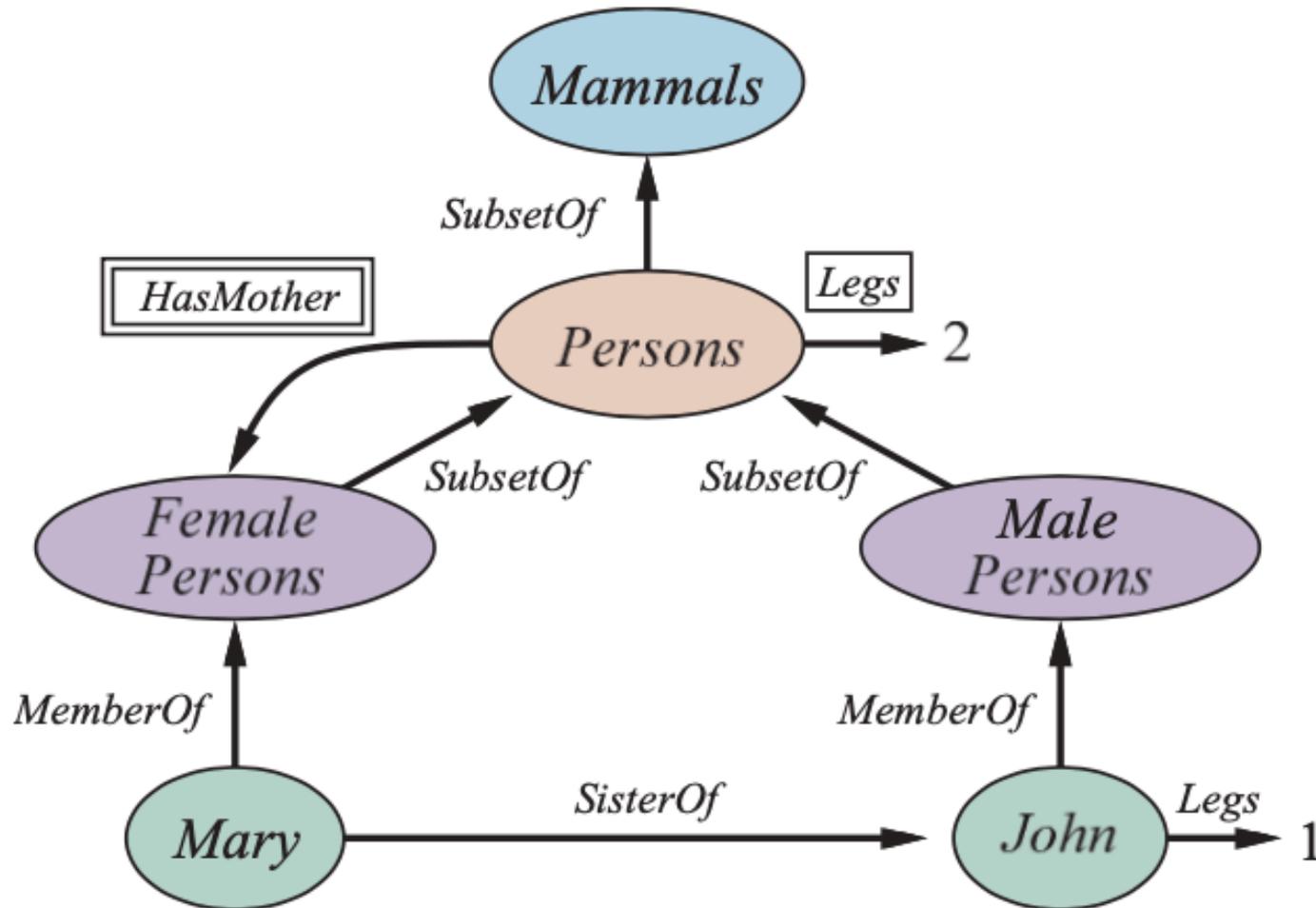


A schematic view of the object President (USA) for the early years



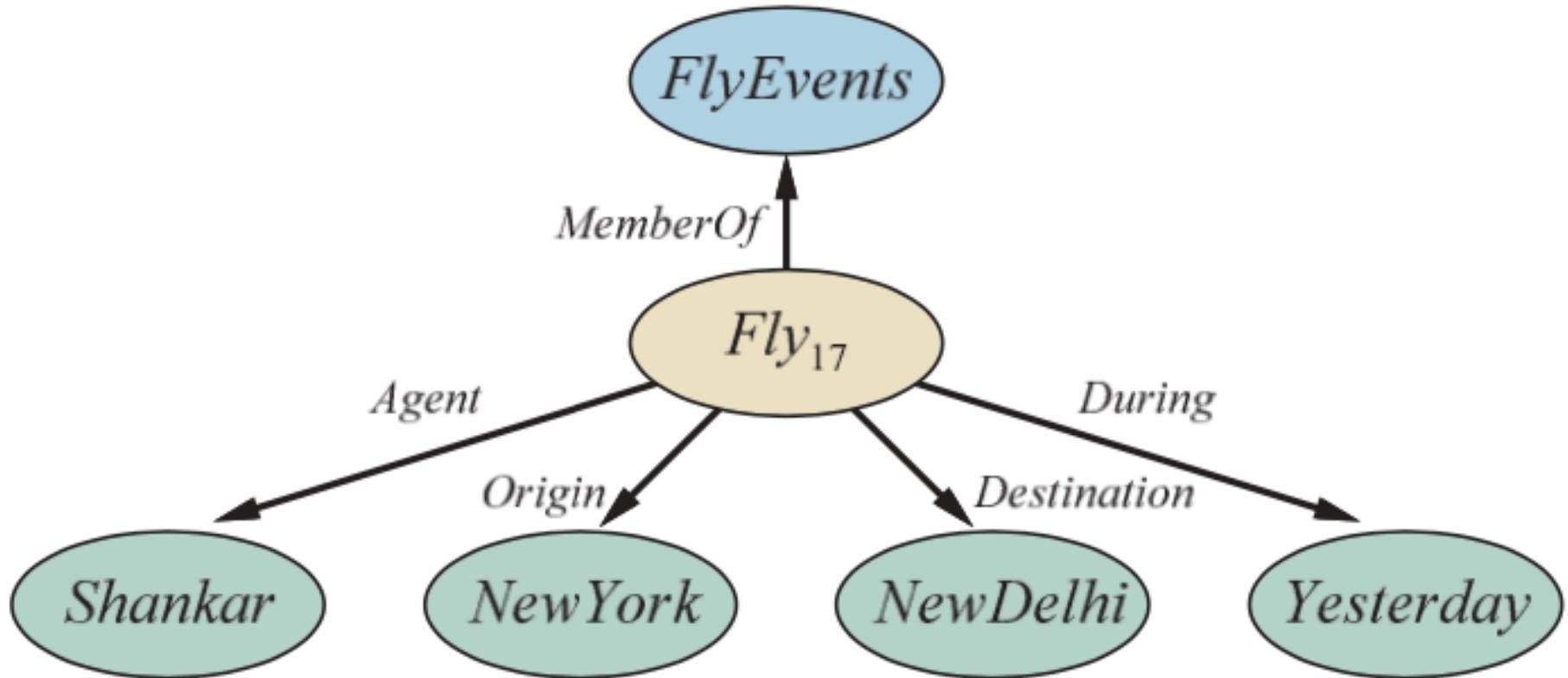
A semantic network

with four objects (John, Mary, 1, and 2) and four categories
Relations are denoted by labeled links



Semantic network

Representation of the logical assertion
Fly (Shankar, NewYork, NewDelhi, Yesterday)



The syntax of descriptions in a subset of the CLASSIC language.

Concept → **Thing** | *ConceptName*
| **And**(*Concept*, ...)
| **All**(*RoleName*, *Concept*)
| **AtLeast**(*Integer*, *RoleName*)
| **AtMost**(*Integer*, *RoleName*)
| **Fills**(*RoleName*, *IndividualName*, ...)
| **SameAs**(*Path*, *Path*)
| **OneOf**(*IndividualName*, ...)

Path → [*RoleName*, ...]

ConceptName → *Adult* | *Female* | *Male* | ...

RoleName → *Spouse* | *Daughter* | *Son* | ...

Automated Planning

A PDDL description of an air cargo transportation planning problem

Init($At(C_1, SFO) \wedge At(C_2, JFK) \wedge At(P_1, SFO) \wedge At(P_2, JFK)$
 $\wedge Cargo(C_1) \wedge Cargo(C_2) \wedge Plane(P_1) \wedge Plane(P_2)$
 $\wedge Airport(JFK) \wedge Airport(SFO)$)

Goal($At(C_1, JFK) \wedge At(C_2, SFO)$)

Action(*Load*(c, p, a),

PRECOND: $At(c, a) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)$

EFFECT: $\neg At(c, a) \wedge In(c, p)$)

Action(*Unload*(c, p, a),

PRECOND: $In(c, p) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)$

EFFECT: $At(c, a) \wedge \neg In(c, p)$)

Action(*Fly*($p, from, to$),

PRECOND: $At(p, from) \wedge Plane(p) \wedge Airport(from) \wedge Airport(to)$

EFFECT: $\neg At(p, from) \wedge At(p, to)$)

The simple spare tire problem

Init(*Tire*(*Flat*) \wedge *Tire*(*Spare*) \wedge *At*(*Flat*, *Axle*) \wedge *At*(*Spare*, *Trunk*))

Goal(*At*(*Spare*, *Axle*))

Action(*Remove*(*obj*, *loc*),

 PRECOND: *At*(*obj*, *loc*)

 EFFECT: \neg *At*(*obj*, *loc*) \wedge *At*(*obj*, *Ground*))

Action(*PutOn*(*t*, *Axle*),

 PRECOND: *Tire*(*t*) \wedge *At*(*t*, *Ground*) \wedge \neg *At*(*Flat*, *Axle*) \wedge \neg *At*(*Spare*, *Axle*)

 EFFECT: \neg *At*(*t*, *Ground*) \wedge *At*(*t*, *Axle*))

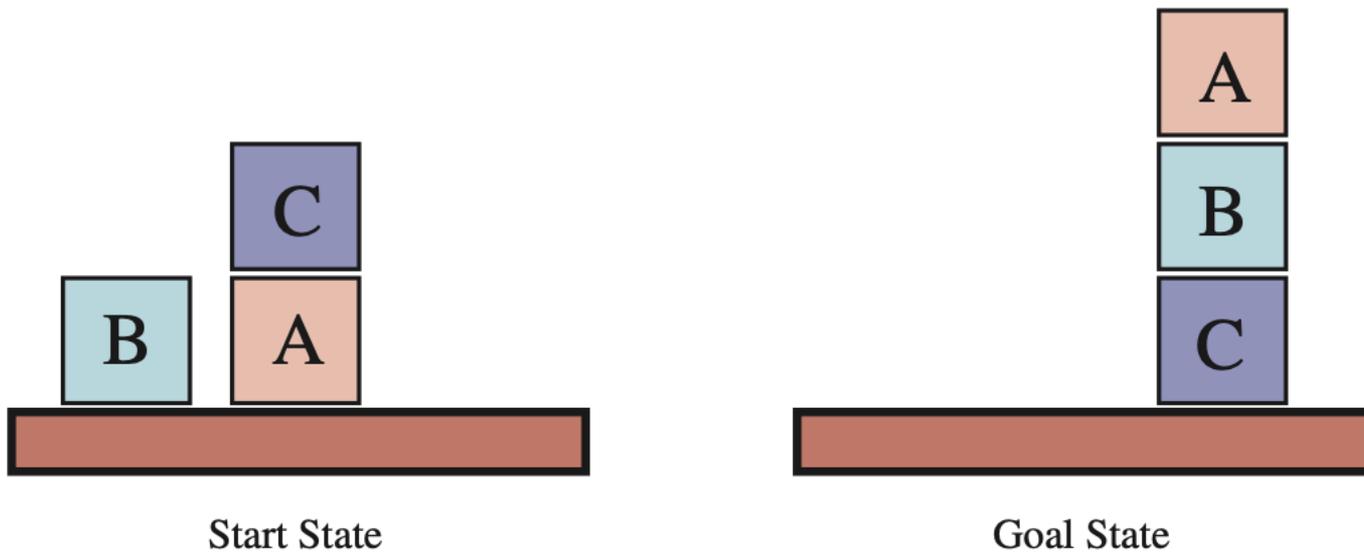
Action(*LeaveOvernight*,

 PRECOND:

 EFFECT: \neg *At*(*Spare*, *Ground*) \wedge \neg *At*(*Spare*, *Axle*) \wedge \neg *At*(*Spare*, *Trunk*)

\wedge \neg *At*(*Flat*, *Ground*) \wedge \neg *At*(*Flat*, *Axle*) \wedge \neg *At*(*Flat*, *Trunk*))

Diagram of the blocks-world problem



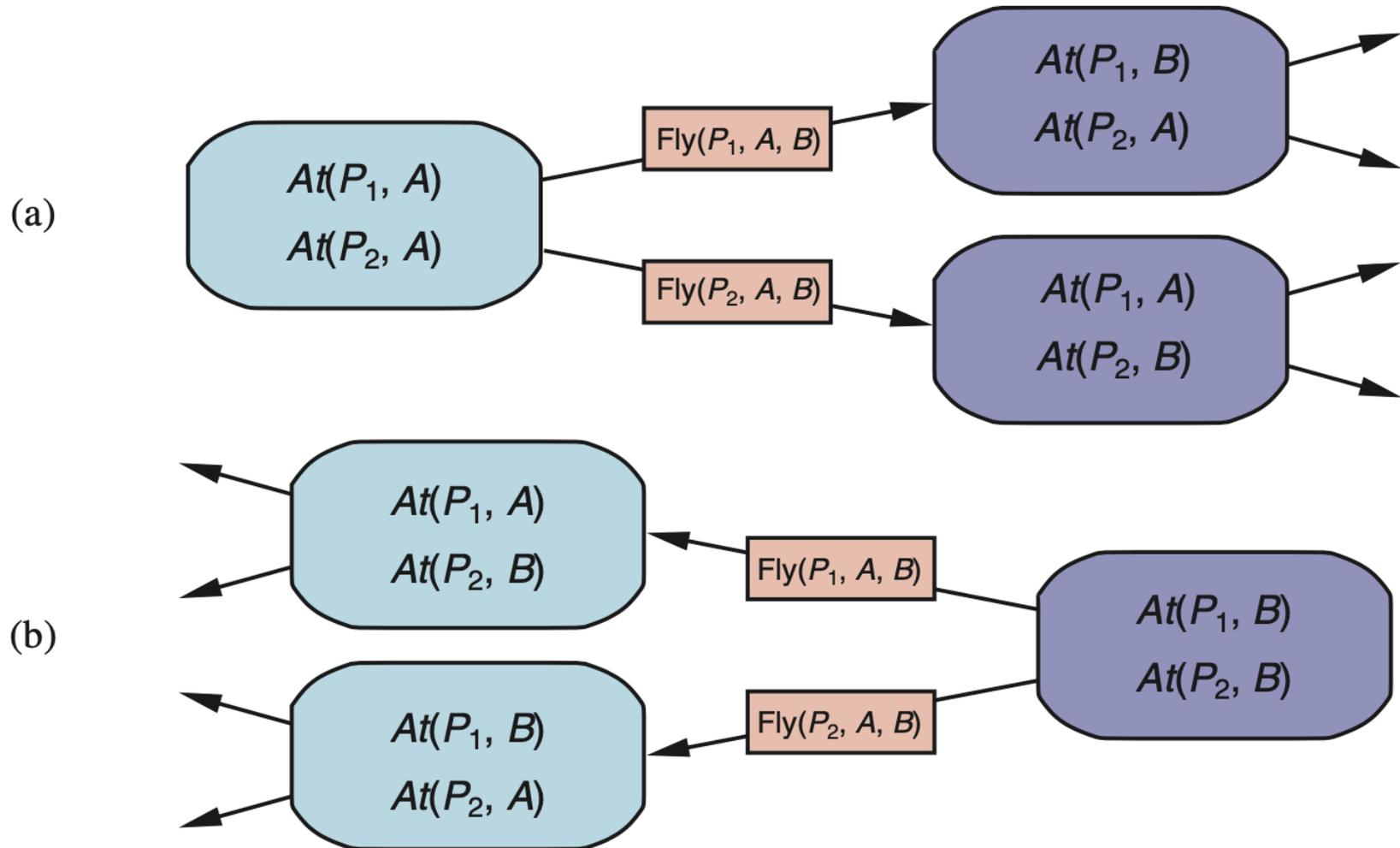
A planning problem in the blocks world: building a three-block tower

Init($On(A, Table) \wedge On(B, Table) \wedge On(C, A)$
 $\wedge Block(A) \wedge Block(B) \wedge Block(C) \wedge Clear(B) \wedge Clear(C) \wedge Clear(Table)$)
Goal($On(A, B) \wedge On(B, C)$)
Action(*Move*(b, x, y),
 PRECOND: $On(b, x) \wedge Clear(b) \wedge Clear(y) \wedge Block(b) \wedge Block(y) \wedge$
 $(b \neq x) \wedge (b \neq y) \wedge (x \neq y)$,
 EFFECT: $On(b, y) \wedge Clear(x) \wedge \neg On(b, x) \wedge \neg Clear(y)$)
Action(*MoveToTable*(b, x),
 PRECOND: $On(b, x) \wedge Clear(b) \wedge Block(b) \wedge Block(x)$,
 EFFECT: $On(b, Table) \wedge Clear(x) \wedge \neg On(b, x)$)

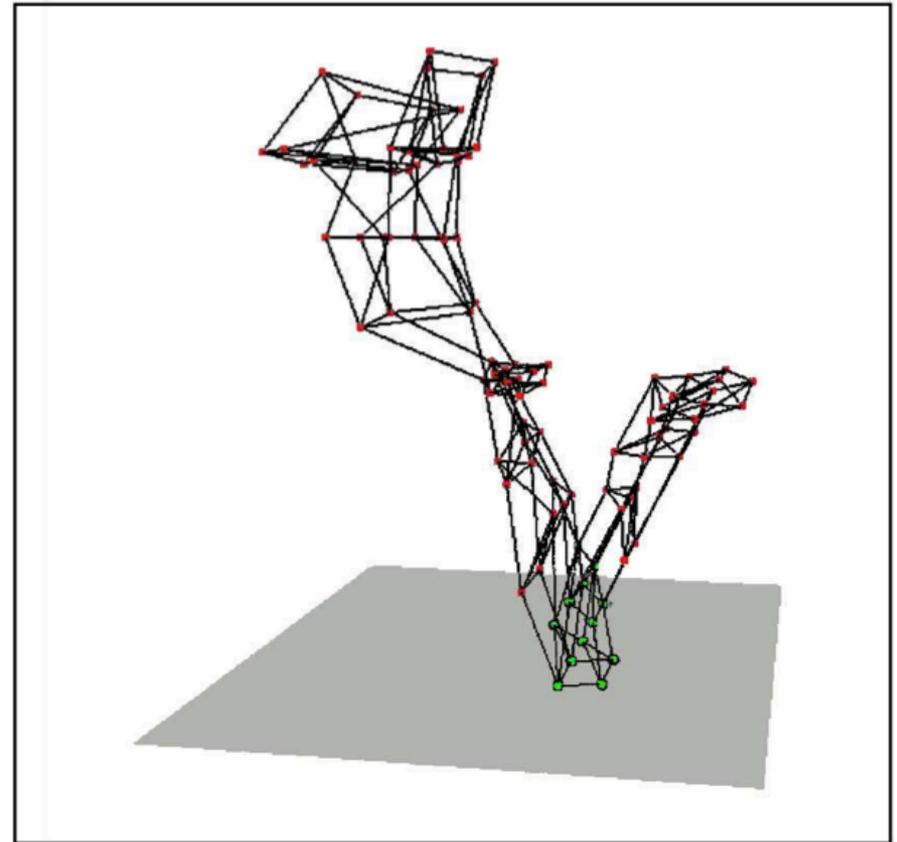
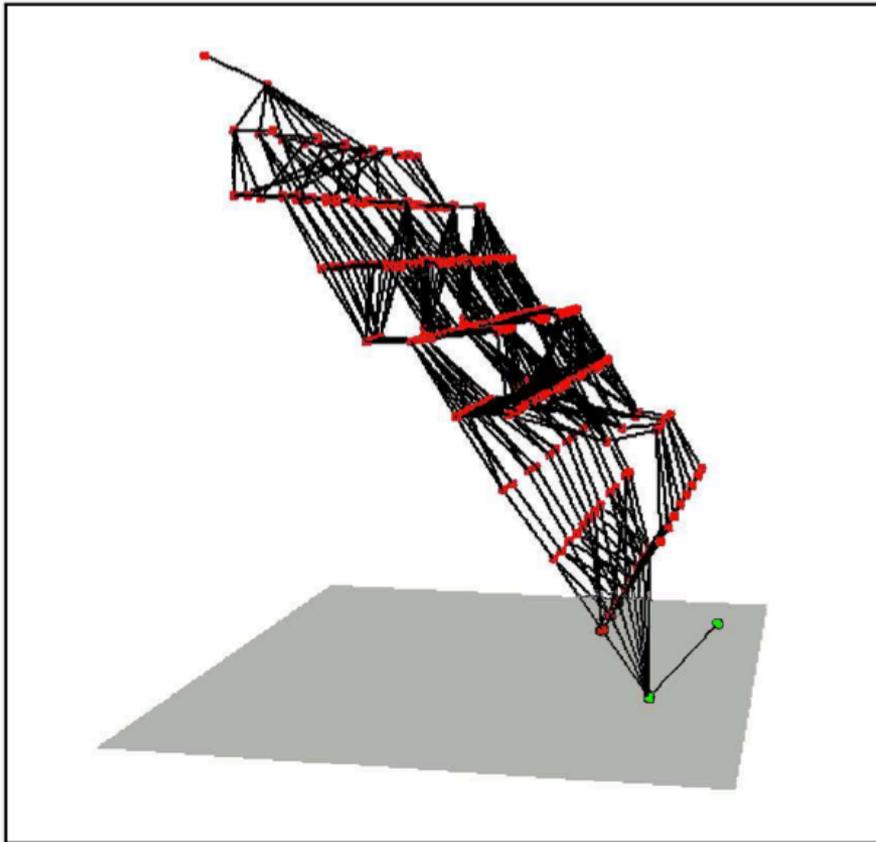
Two approaches to searching for a plan

(a) Forward (progression) search

(b) Backward (regression) search



Two state spaces from planning problems with the ignore-delete-lists heuristic



Definitions of possible refinements for two high-level actions

Refinement(*Go*(*Home*, *SFO*),
 STEPS: [*Drive*(*Home*, *SFO* *LongTermParking*),
 Shuttle(*SFO* *LongTermParking*, *SFO*)])

Refinement(*Go*(*Home*, *SFO*),
 STEPS: [*Taxi*(*Home*, *SFO*)])

Refinement(*Navigate*($[a, b]$, $[x, y]$),
 PRECOND: $a = x \wedge b = y$
 STEPS: [])

Refinement(*Navigate*($[a, b]$, $[x, y]$),
 PRECOND: *Connected*($[a, b]$, $[a - 1, b]$)
 STEPS: [*Left*, *Navigate*($[a - 1, b]$, $[x, y]$)])

Refinement(*Navigate*($[a, b]$, $[x, y]$),
 PRECOND: *Connected*($[a, b]$, $[a + 1, b]$)
 STEPS: [*Right*, *Navigate*($[a + 1, b]$, $[x, y]$)])

A breadth-first implementation of hierarchical forward planning search

function HIERARCHICAL-SEARCH(*problem*, *hierarchy*) **returns** a solution or *failure*

frontier \leftarrow a FIFO queue with [*Act*] as the only element

while *true* **do**

if IS-EMPTY(*frontier*) **then return** *failure*

plan \leftarrow POP(*frontier*) // chooses the shallowest plan in frontier

hla \leftarrow the first HLA in *plan*, or *null* if none

prefix, suffix \leftarrow the action subsequences before and after *hla* in *plan*

outcome \leftarrow RESULT(*problem*.INITIAL, *prefix*)

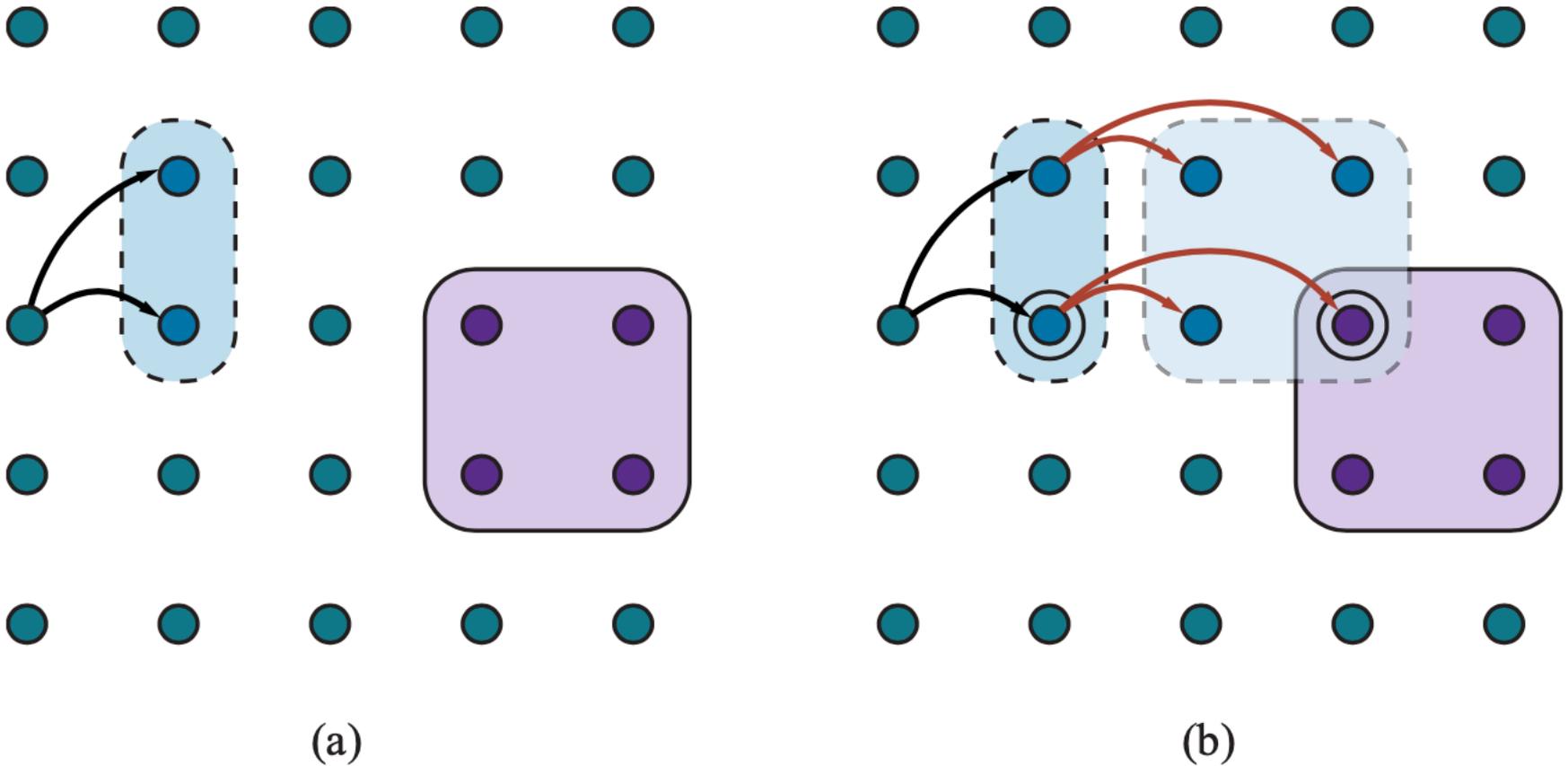
if *hla* is *null* **then** // so *plan* is primitive and *outcome* is its result

if *problem*.IS-GOAL(*outcome*) **then return** *plan*

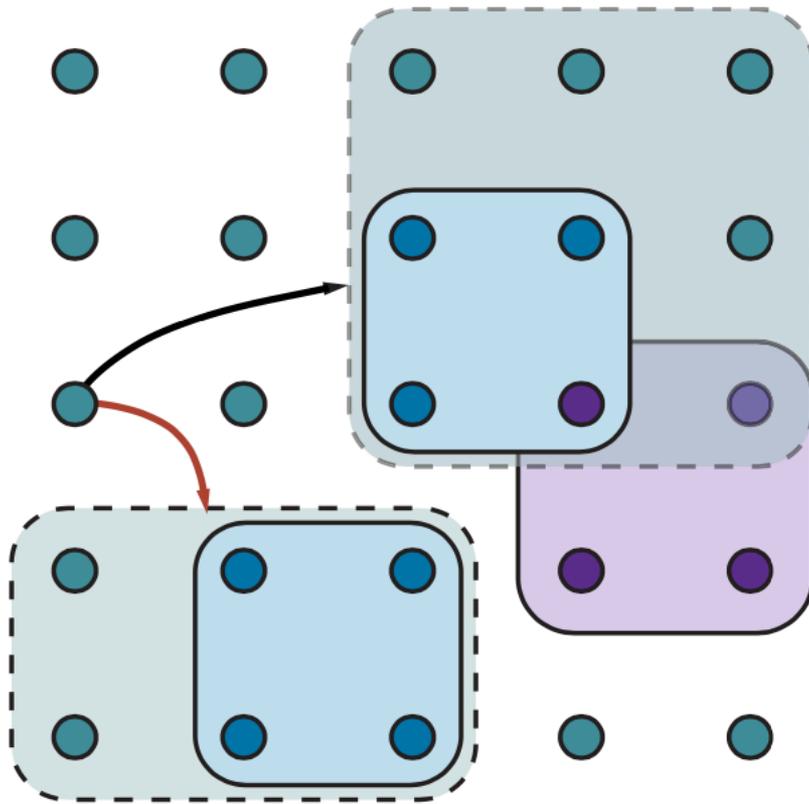
else for each *sequence* **in** REFINEMENTS(*hla*, *outcome*, *hierarchy*) **do**

 add APPEND(*prefix*, *sequence*, *suffix*) to *frontier*

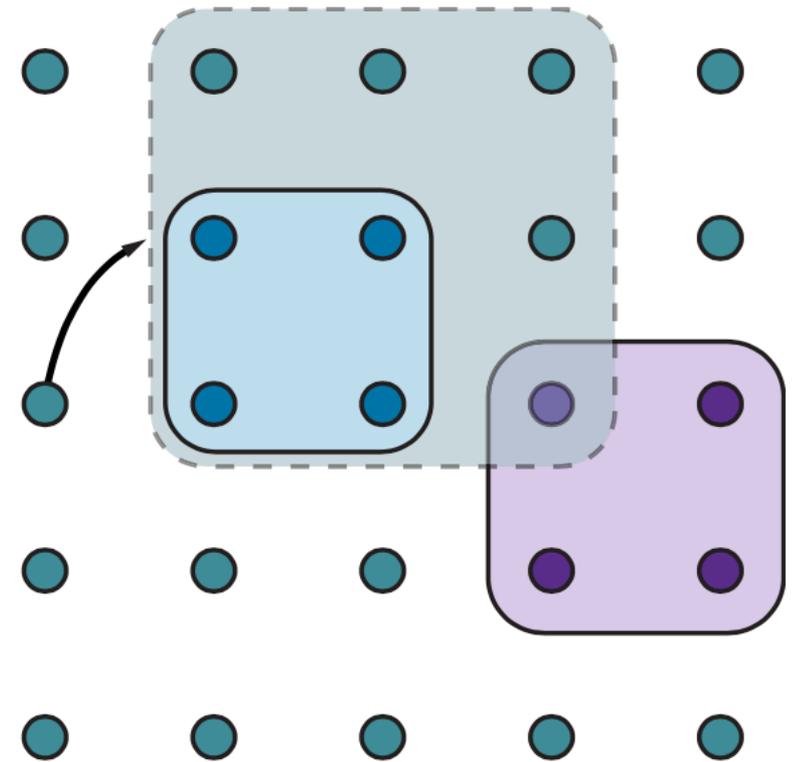
Schematic examples of reachable sets



Goal achievement for high-level plans with approximate descriptions



(a)



(b)

A hierarchical planning algorithm

```
function ANGELIC-SEARCH(problem, hierarchy, initialPlan) returns solution or fail  
frontier  $\leftarrow$  a FIFO queue with initialPlan as the only element  
while true do  
  if EMPTY?(frontier) then return fail  
  plan  $\leftarrow$  POP(frontier) // chooses the shallowest node in frontier  
  if REACH+(problem.INITIAL, plan) intersects problem.GOAL then  
    if plan is primitive then return plan // REACH+ is exact for primitive plans  
    guaranteed  $\leftarrow$  REACH-(problem.INITIAL, plan)  $\cap$  problem.GOAL  
    if guaranteed  $\neq$  { } and MAKING-PROGRESS(plan, initialPlan) then  
      finalState  $\leftarrow$  any element of guaranteed  
      return DECOMPOSE(hierarchy, problem.INITIAL, plan, finalState)  
    hla  $\leftarrow$  some HLA in plan  
    prefix, suffix  $\leftarrow$  the action subsequences before and after hla in plan  
    outcome  $\leftarrow$  RESULT(problem.INITIAL, prefix)  
    for each sequence in REFINEMENTS(hla, outcome, hierarchy) do  
      frontier  $\leftarrow$  Insert(APPEND(prefix, sequence, suffix), frontier)
```

A hierarchical planning algorithm

Decompose solution

function DECOMPOSE(*hierarchy*, s_0 , *plan*, s_f) **returns** a solution

solution \leftarrow an empty plan

while *plan* is not empty **do**

action \leftarrow REMOVE-LAST(*plan*)

s_i \leftarrow a state in $\text{REACH}^-(s_0, \text{plan})$ such that $s_f \in \text{REACH}^-(s_i, \text{action})$

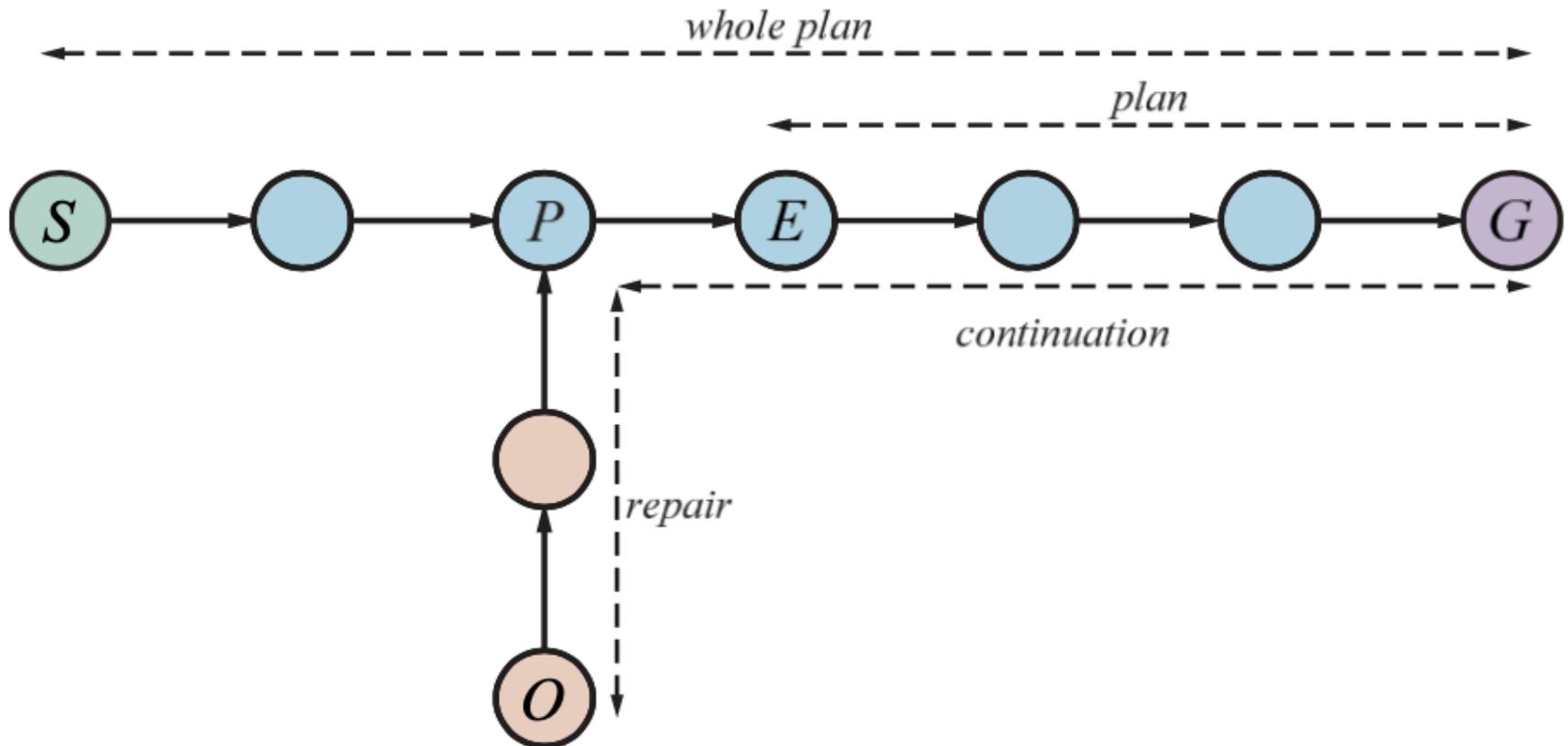
problem \leftarrow a problem with INITIAL = s_i and GOAL = s_f

solution \leftarrow APPEND(ANGELIC-SEARCH(*problem*, *hierarchy*, *action*), *solution*)

s_f \leftarrow s_i

return *solution*

At first, the sequence “whole plan”
is expected to
get the agent from S to G



A job-shop scheduling problem for assembling two cars, with resource constraints

Jobs({*AddEngine1* \prec *AddWheels1* \prec *Inspect1* },
{*AddEngine2* \prec *AddWheels2* \prec *Inspect2* })

Resources(*EngineHoists*(1), *WheelStations*(1), *Inspectors*(e2), *LugNuts*(500))

Action(*AddEngine1*, DURATION:30,
USE:*EngineHoists*(1))

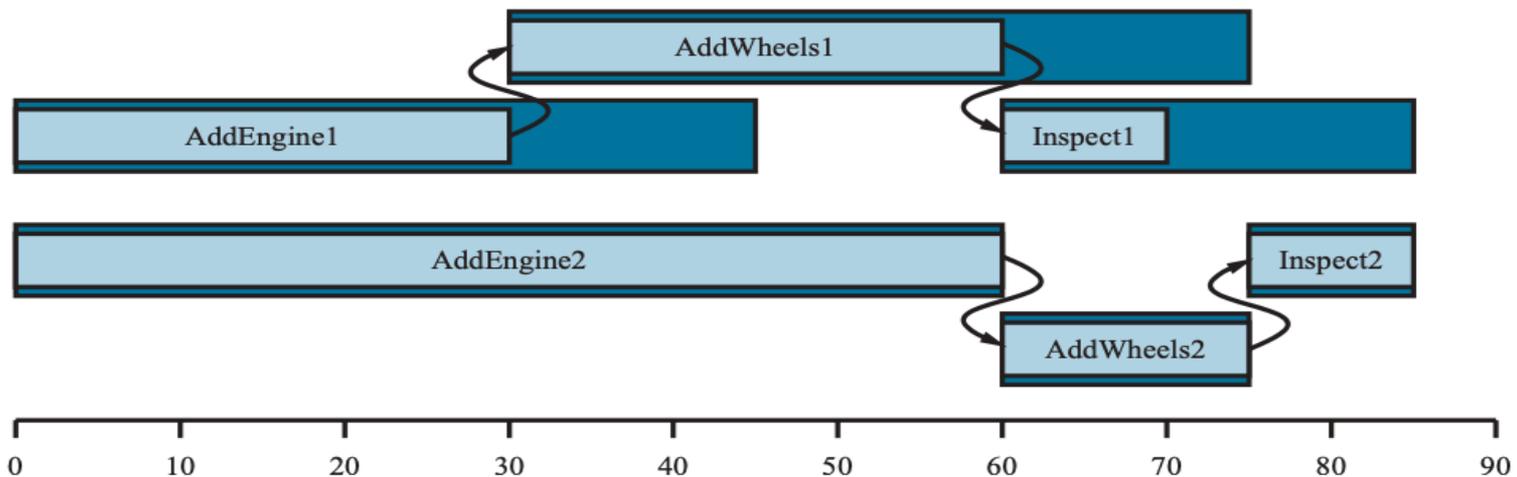
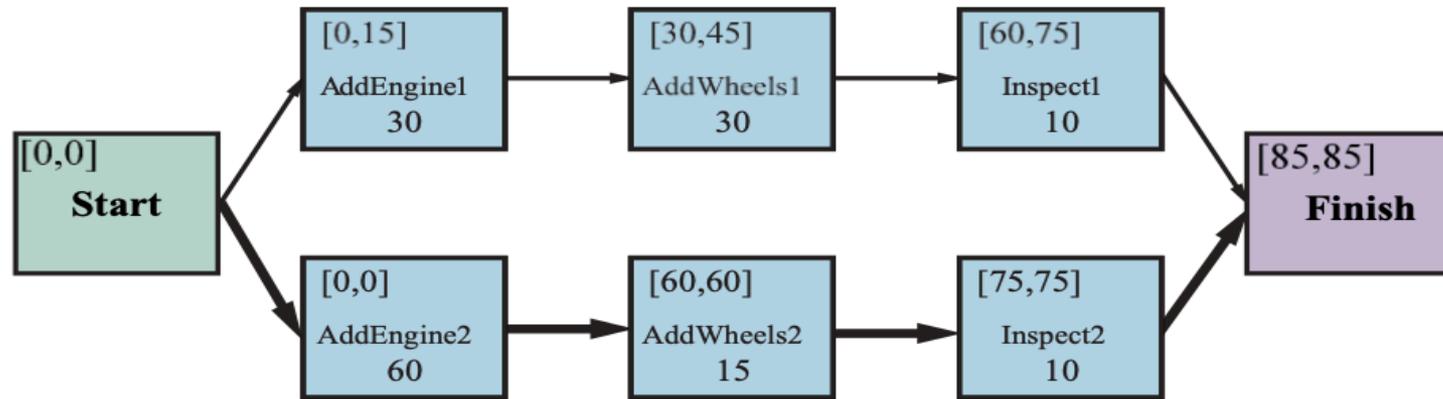
Action(*AddEngine2*, DURATION:60,
USE:*EngineHoists*(1))

Action(*AddWheels1*, DURATION:30,
CONSUME:*LugNuts*(20), USE:*WheelStations*(1))

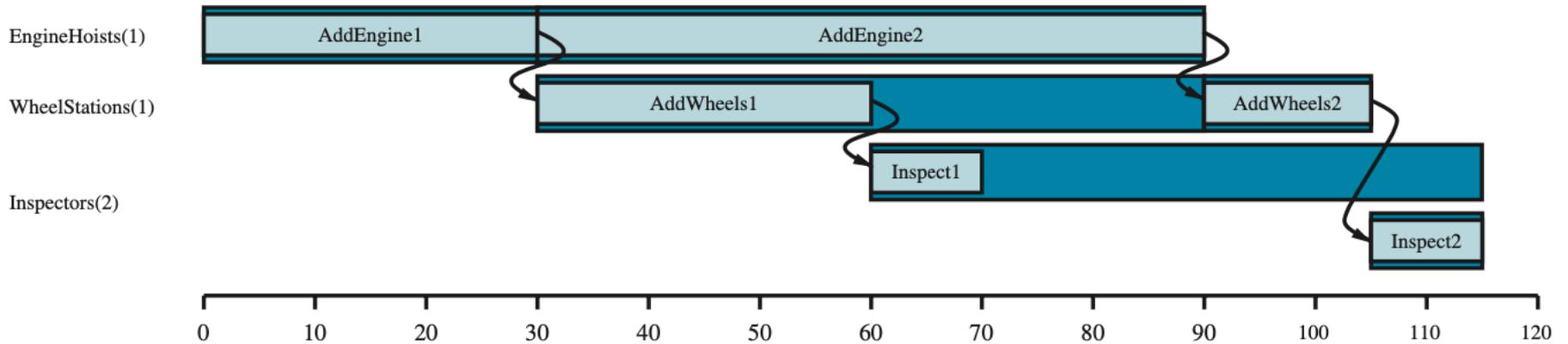
Action(*AddWheels2*, DURATION:15,
CONSUME:*LugNuts*(20), USE:*WheelStations*(1))

Action(*Inspect_i*, DURATION:10,
USE:*Inspectors*(1))

A representation of the temporal constraints for the job-shop scheduling problem



A solution to the job-shop scheduling problem

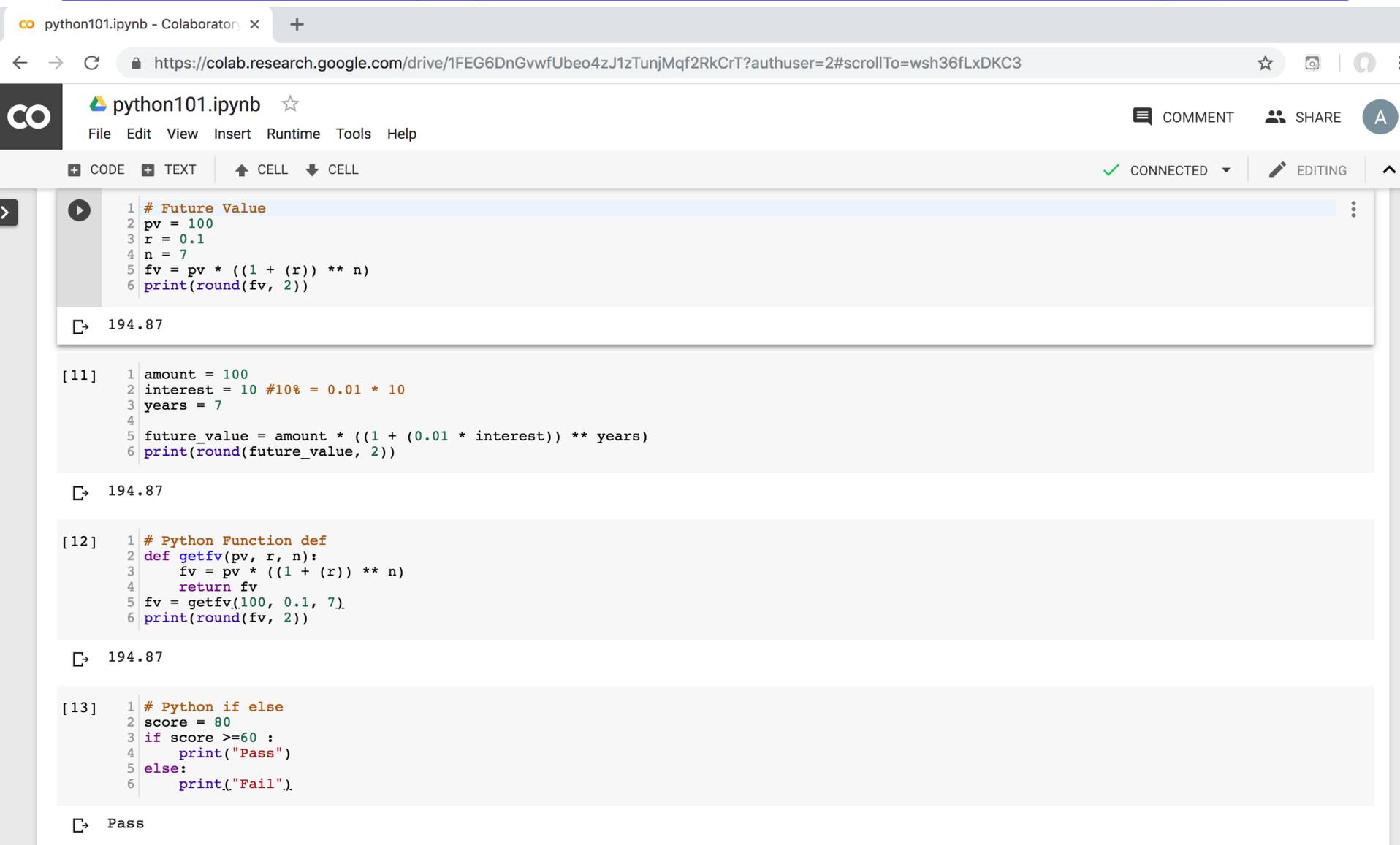


AIMA Python

- Artificial Intelligence: A Modern Approach (AIMA)
 - <http://aima.cs.berkeley.edu/>
- AIMA Python
 - <http://aima.cs.berkeley.edu/python/readme.html>
- Logic, KB Agent
 - <http://aima.cs.berkeley.edu/python/logic.html>

Python in Google Colab (Python101)

<https://colab.research.google.com/drive/1FEG6DnGvwfUbeo4zJ1zTunjMqf2RkCrT>



python101.ipynb - Collaborator

File Edit View Insert Runtime Tools Help

COMMENT SHARE

CONNECTED EDITING

```
1 # Future Value
2 pv = 100
3 r = 0.1
4 n = 7
5 fv = pv * ((1 + (r)) ** n)
6 print(round(fv, 2))
```

194.87

```
[11] 1 amount = 100
2 interest = 10 #10% = 0.01 * 10
3 years = 7
4
5 future_value = amount * ((1 + (0.01 * interest)) ** years)
6 print(round(future_value, 2))
```

194.87

```
[12] 1 # Python Function def
2 def getfv(pv, r, n):
3     fv = pv * ((1 + (r)) ** n)
4     return fv
5 fv = getfv(100, 0.1, 7)
6 print(round(fv, 2))
```

194.87

```
[13] 1 # Python if else
2 score = 80
3 if score >=60 :
4     print("Pass")
5 else:
6     print("Fail").
```

Pass

<https://tinyurl.com/aintpupython101>

Summary

- **Logical Agents**
- **First-Order Logic**
- **Inference in First-Order Logic**
- **Knowledge Representation**
- **Automated Planning**

References

- Stuart Russell and Peter Norvig (2020), Artificial Intelligence: A Modern Approach, 4th Edition, Pearson.
- Aurélien Géron (2019), Hands-On Machine Learning with Scikit-Learn, Keras, and TensorFlow: Concepts, Tools, and Techniques to Build Intelligent Systems, 2nd Edition, O'Reilly Media.