## Artificial Intelligence

## Knowledge，Reasoning and Knowledge Representation

## Uncertain Knowledge and Reasoning

1111AI04<br>MBA，IM，NTPU（M6132）（Fall 2022）<br>Wed 2，3， 4 （9：10－12：00）（B8F40）



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## Syllabus

## Week Date Subject／Topics

1 2022／09／14 Introduction to Artificial Intelligence
2 2022／09／21 Artificial Intelligence and Intelligent Agents
3 2022／09／28 Problem Solving
4 2022／10／05 Knowledge，Reasoning and Knowledge Representation； Uncertain Knowledge and Reasoning

5 2022／10／12 Case Study on Artificial Intelligence I
6 2022／10／19 Machine Learning：Supervised and Unsupervised Learning

## Syllabus

Week Date Subject／Topics
7 2022／10／26 The Theory of Learning and Ensemble Learning
8 2022／11／02 Midterm Project Report
9 2022／11／09 Deep Learning and Reinforcement Learning
10 2022／11／16 Deep Learning for Natural Language Processing
11 2022／11／23 Invited Talk：AI for Information Retrieval
12 2022／11／30 Case Study on Artificial Intelligence II

## Syllabus

Week Date Subject/Topics
13 2022/12/07 Computer Vision and Robotics
14 2022/12/14 Philosophy and Ethics of AI and the Future of AI
15 2022/12/21 Final Project Report I
16 2022/12/28 Final Project Report II
17 2023/01/04 Self-learning
18 2023/01/11 Self-learning

# Knowledge, Reasoning and Knowledge Representation 

## Uncertain Knowledge and Reasoning

## Outline

- Knowledge and Reasoning
- Logical Agents
- First-Order Logic
- Inference in First-Order Logic
- Knowledge Representation
- Knowledge Graph (KG)
- Uncertain Knowledge and Reasoning
- Quantifying Uncertainty
- Probabilistic Reasoning
- Making Complex Decisions

Stuart Russell and Peter Norvig (2020), Artificial Intelligence: A Modern Approach,

4th Edition, Pearson


## Artificial Intelligence: A Modern Approach

1. Artificial Intelligence
2. Problem Solving
3. Knowledge and Reasoning
4. Uncertain Knowledge and Reasoning

## 5. Machine Learning

6. Communicating, Perceiving, and Acting
7. Philosophy and Ethics of AI

# Artificial Intelligence: 

 Knowledge and Reasoning
# Artificial Intelligence: 3. Knowledge and Reasoning 

- Logical Agents
- First-Order Logic
- Inference in First-Order Logic
- Knowledge Representation
- Automated Planning


## Intelligent Agents

## 4 Approaches of AI

$\left.\begin{array}{|c|c|}\hline \text { 2. } & \begin{array}{c}3 . \\ \text { Thinking Humanly: } \\ \text { The Cognitive } \\ \text { Modeling Approach }\end{array}\end{array} \begin{array}{c}\text { Thinking Rationally: } \\ \text { The "Laws of Thought" } \\ \text { Approach }\end{array}\right]$

## Reinforcement Learning (DL)

## Agent

## Environment

## Reinforcement Learning (DL)



## Reinforcement Learning (DL)



## Agents interact with environments through sensors and actuators



## Logical Agents

## Logical Agents

## Knowledge-based Agents KB Agents

## Knowledge-based Agent (KB Agent)

function KB-AGENT( percept) returns an action persistent: $K B$, a knowledge base $t$, a counter, initially 0 , indicating time

Tell( $K B$, Make-Percept-Sentence percept, $t$ )) action $\leftarrow \operatorname{AsK}(K B$, MAKE-ACtion-Query $(t))$ Tell( $K B$, MaKe-Action-Sentence $($ action, $t)$ ) $t \leftarrow t+1$

return action

## Sentences are physical configurations of the agent



Reasoning is a process of constructing new physical configurations from old ones

Logical reasoning should ensure that the new configurations represent aspects of the world that actually follow from the aspects that the old configurations represent.

## A BNF (Backus-Naur Form) grammar of sentences in propositional logic

Sentence $\rightarrow$ AtomicSentence $\mid$ ComplexSentence<br>AtomicSentence $\rightarrow$ True $\mid$ False $|P| Q|R| \ldots$<br>ComplexSentence $\rightarrow$ (Sentence )<br>| $\neg$ Sentence<br>| Sentence $\wedge$ Sentence<br>Sentence $\vee$ Sentence<br>Sentence $\Rightarrow$ Sentence<br>| Sentence $\Leftrightarrow$ Sentence

Operator Precedence $\quad: \quad \neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$

## Truth Tables (TT) for the Five Logical Connectives

| $P$ | $Q$ | $\neg P$ | $P \wedge Q$ | $P \vee Q$ | $P \Rightarrow Q$ | $P \Leftrightarrow Q$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| false | false | true | false | false | true | true |
| false | true | true | false | true | true | false |
| true | false | false | false | true | false | false |
| true | true | false | true | true | true | true |

## A Truth Table constructed for the knowledge base given in the text

| $B_{1,1}$ | $B_{2,1}$ | $P_{1,1}$ | $P_{1,2}$ | $P_{2,1}$ | $P_{2,2}$ | $P_{3,1}$ | $R_{1}$ | $R_{2}$ | $R_{3}$ | $R_{4}$ | $R_{5}$ | KB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| false | false | false | false | false | false | false | true | true | true | true | false | false |
| false | false | false | false | false | false | true | true | true | false | true | false | false |
| $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ |
| false | true | false | false | false | false | false | true | true | false | true | true | false |
| false | true | false | false | false | false | true | true | true | true | true | true | true |
| false | true | false | false | false | true | false | true | true | true | true | true | $\underline{\text { true }}$ |
| false | true | false | false | false | true | true | true | true | true | true | true | true |
| false | true | false | false | true | false | false | true | false | false | true | true | false |
| $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ |
| true | true | true | true | true | true | true | false | true | true | false | true | false |

# A Truth-Table (TT) enumeration algorithm for deciding propositional entailment 

```
function TT-EnTAILS?(KB,\alpha) returns true or false
    inputs: KB, the knowledge base, a sentence in propositional logic
            \alpha, the query, a sentence in propositional logic
    symbols }\leftarrow\textrm{a}\mathrm{ list of the proposition symbols in KB and }
    return TT-CHECK-ALL(KB, \alpha, symbols, { })
function TT-CHECK-AlL(KB, \alpha, symbols, model) returns true or false
    if EmpTY?(symbols) then
    if PL-TruE?(KB, model) then return PL-TruE?( }\alpha,\mathrm{ model)
    else return true // when KB is false, always return true
    else
        P\leftarrowFIRST(symbols)
        rest }\leftarrow\mathrm{ REST(symbols)
        return(TT-CHECK-ALL(KB, \alpha, rest, model \cup{P=true})
            and
            TT-CHECK-ALL(KB, }\alpha\mathrm{ , rest, model }\cup{P=\mathrm{ false }))
```


## Standard Logical Equivalences

The symbols $\alpha, \beta$, and $\gamma$ stand for arbitrary sentences of propositional logic.

$$
\begin{aligned}
(\alpha \wedge \beta) & \equiv(\beta \wedge \alpha) \quad \text { commutativity of } \wedge \\
(\alpha \vee \beta) & \equiv(\beta \vee \alpha) \quad \text { commutativity of } \vee \\
((\alpha \wedge \beta) \wedge \gamma) & \equiv(\alpha \wedge(\beta \wedge \gamma)) \quad \text { associativity of } \wedge \\
((\alpha \vee \beta) \vee \gamma) & \equiv(\alpha \vee(\beta \vee \gamma)) \quad \text { associativity of } \vee \\
\neg(\neg \alpha) & \equiv \alpha \text { double-negation elimination } \\
(\alpha \Rightarrow \beta) & \equiv(\neg \beta \Rightarrow \neg \alpha) \quad \text { contraposition } \\
(\alpha \Rightarrow \beta) & \equiv(\neg \alpha \vee \beta) \quad \text { implication elimination } \\
(\alpha \Leftrightarrow \beta) & \equiv((\alpha \Rightarrow \beta) \wedge(\beta \Rightarrow \alpha)) \quad \text { biconditional elimination } \\
\neg(\alpha \wedge \beta) & \equiv(\neg \alpha \vee \neg \beta) \quad \text { De Morgan } \\
\neg(\alpha \vee \beta) & \equiv(\neg \alpha \wedge \neg \beta) \quad \text { De Morgan } \\
(\alpha \wedge(\beta \vee \gamma)) & \equiv((\alpha \wedge \beta) \vee(\alpha \wedge \gamma)) \text { distributivity of } \wedge \text { over } \vee \\
(\alpha \vee(\beta \wedge \gamma)) & \equiv((\alpha \vee \beta) \wedge(\alpha \vee \gamma)) \quad \text { distributivity of } \vee \text { over } \wedge
\end{aligned}
$$

# A grammar for Conjunctive Normal Form (CNF), Horn clauses, and definite clauses 

$$
\begin{aligned}
& \text { CNFSentence } \rightarrow \text { Clause }_{1} \wedge \cdots \wedge \text { Clause }_{n} \\
& \text { Clause } \rightarrow \text { Literal }_{1} \vee \cdots \vee \text { Literal }_{m} \\
& \text { Fact } \rightarrow \text { Symbol } \\
& \text { Literal } \rightarrow \text { Symbol } \mid \neg \text { Symbol } \\
& \text { Symbol } \rightarrow P|Q| R \mid \ldots \\
& \text { HornClauseForm } \rightarrow \text { DefiniteClauseForm } \mid \text { GoalClauseForm } \\
& \text { DefiniteClauseForm } \rightarrow \text { Fact } \mid\left(\text { Symbol }_{1} \wedge \cdots \wedge \text { Symbol }_{l}\right) \Rightarrow \text { Symbol } \\
& \text { GoalClauseForm } \rightarrow\left(\text { Symbol }_{1} \wedge \cdots \wedge \text { Symbol }_{l}\right) \Rightarrow \text { False }
\end{aligned}
$$

## A simple resolution algorithm for propositional logic

function PL-RESOLUTION $(K B, \alpha)$ returns true or false inputs: $K B$, the knowledge base, a sentence in propositional logic $\alpha$, the query, a sentence in propositional logic
clauses $\leftarrow$ the set of clauses in the CNF representation of $K B \wedge \neg \alpha$ new $\leftarrow\}$
while true do
for each pair of clauses $C_{i}, C_{j}$ in clauses do resolvents $\leftarrow \operatorname{PL}-\operatorname{ReSOLVE}\left(C_{i}, C_{j}\right)$
if resolvents contains the empty clause then return true new $\leftarrow$ new $\cup$ resolvents
if new $\subseteq$ clauses then return false clauses $\leftarrow$ clauses $\cup$ new

## The forward-chaining algorithm for propositional logic

function PL-FC-Entails? $(K B, q)$ returns true or false
inputs: $K B$, the knowledge base, a set of propositional definite clauses
$q$, the query, a proposition symbol
count $\leftarrow$ a table, where count $[c]$ is initially the number of symbols in clause $c$ 's premise inferred $\leftarrow$ a table, where inferred $[s]$ is initially false for all symbols queue $\leftarrow$ a queue of symbols, initially symbols known to be true in $K B$
while queue is not empty do
$p \leftarrow \operatorname{POP}(q u e u e)$
if $p=q$ then return true
if inferred $[p]=$ false then
inferred $[p] \leftarrow$ true
for each clause $c$ in $K B$ where $p$ is in $c$. PREMISE do
decrement count $[c]$
if count $[c]=0$ then add $c$. CONCLUSION to queue
return false

## A set of Horn clauses

$$
\begin{aligned}
& P \Rightarrow Q \\
& L \wedge M \Rightarrow P \\
& B \wedge L \Rightarrow M \\
& A \wedge P \Rightarrow L \\
& A \wedge B \Rightarrow L \\
& A \\
& B
\end{aligned}
$$


(b)

The corresponding AND-OR graph

## First-Order Logic

## Formal languages and their ontological and epistemological commitments

| Language | Ontological Commitment <br> (What exists in the world) | Epistemological Commitment <br> (What an agent believes about facts) |
| :--- | :--- | :--- |
| Propositional logic | facts | true/false/unknown |
| First-order logic | facts, objects, relations | true/false/unknown |
| Temporal logic | facts, objects, relations, times | true/false/unknown |
| Probability theory | facts | degree of belief $\in[0,1]$ |
| Fuzzy logic | facts with degree of truth $\in[0,1]$ | known interval value |

## A model containing five objects

two binary relations (brother and on-head), three unary relations (person, king, and crown), and one unary function (left-leg).


The syntax of first-order logic with equality


Some members of the set of all models for a language with two constant symbols, R and J, and one binary relation symbol


Some members of the set of all models for a language with two constant symbols, $R$ and J, and one binary relation symbol, under database semantics


A digital circuit C1, purporting to be a one-bit full adder.


# Inference in First-Order Logic 

## The unification algorithm

function $\operatorname{UNIFY}(x, y, \theta=$ empty $)$ returns a substitution to make $x$ and $y$ identical, or failure
if $\theta=$ failure then return failure
else if $x=y$ then return $\theta$
else if $\operatorname{VARIABLE} ?(x)$ then return $\operatorname{Unify-VAR}(x, y, \theta)$
else if Variable? $(y)$ then return $\operatorname{Unify-\operatorname {Var}(y,x,\theta )}$
else if Compound? $(x)$ and Compound? $(y)$ then
return $\operatorname{Unify}(\operatorname{ArgS}(x), \operatorname{ArgS}(y), \operatorname{Unify}(\operatorname{OP}(x), \operatorname{OP}(y), \theta))$
else if LISt? $(x)$ and List? ( $y$ ) then
return $\operatorname{UnIFY}(\operatorname{Rest}(x), \operatorname{Rest}(y), \operatorname{UnIFY}(\operatorname{First}(x), \operatorname{FIRST}(y), \theta))$
else return failure
function UNIFY-VAR $(v a r, x, \theta)$ returns a substitution
if $\{$ var $/$ val $\} \in \theta$ for some val then return $\operatorname{UNIFY}($ val, $x, \theta)$
else if $\{x /$ val $\} \in \theta$ for some val then return UNify (var, val, $\theta$ )
else if OCCUR-CHECK? $(\mathrm{var}, x)$ then return failure
else return add $\{\operatorname{var} / x\}$ to $\theta$

## The subsumption lattice whose lowest node is Employs (IBM , Richard)

The subsumption lattice for the sentence Employs (John, John )

(a)

(b)

## A conceptually straightforward, but inefficient, forward-chaining algorithm

function FOL-FC-ASK $(K B, \alpha)$ returns a substitution or false
inputs: $K B$, the knowledge base, a set of first-order definite clauses
$\alpha$, the query, an atomic sentence
while true do
new $\leftarrow\} \quad / /$ The set of new sentences inferred on each iteration
for each rule in $K B$ do
$\left(p_{1} \wedge \ldots \wedge p_{n} \Rightarrow q\right) \leftarrow$ Standardize-VARIABLES $($ rule $)$
for each $\theta$ such that $\operatorname{SUBST}\left(\theta, p_{1} \wedge \ldots \wedge p_{n}\right)=\operatorname{Subst}\left(\theta, p_{1}^{\prime} \wedge \ldots \wedge p_{n}^{\prime}\right)$
for some $p_{1}^{\prime}, \ldots, p_{n}^{\prime}$ in $K B$
$q^{\prime} \leftarrow \operatorname{SUBST}(\theta, q)$
if $q^{\prime}$ does not unify with some sentence already in $K B$ or new then add $q^{\prime}$ to new $\phi \leftarrow \operatorname{UNiFY}\left(q^{\prime}, \alpha\right)$ if $\phi$ is not failure then return $\phi$
if $n e w=\{ \}$ then return false
add new to $K B$

## The proof tree generated by forward chaining on the crime example



## Constraint graph for coloring the map of Australia


(a)

$$
\begin{aligned}
& \text { Diff }(w a, n t) \wedge \text { Diff }(w a, s a) \wedge \\
& \text { Diff }(n t, q) \wedge \text { Diff }(n t, s a) \wedge \\
& \text { Diff }(q, n s w) \wedge \text { Diff }(q, s a) \wedge \\
& \text { Diff }(n s w, v) \wedge \text { Diff }(n s w, \text { sa }) \wedge \\
& \text { Diff }(v, \text { sa }) \Rightarrow \text { Colorable }() \\
& \text { Diff }(\text { Red, Blue }) \quad \text { Diff }(\text { Red, Green }) \\
& \text { Diff }(\text { Green }, \text { Red }) \text { Diff }(\text { Green , Blue }) \\
& \text { Diff }(\text { Blue }, \text { Red }) \quad \text { Diff (Blue, Green })
\end{aligned}
$$

(b)

## A simple backward-chaining algorithm for first-order knowledge bases

function FOL-BC-ASK (KB, query) returns a generator of substitutions return $\mathrm{FOL}-\mathrm{BC}-\mathrm{OR}(K B$, query, $\{ \})$
function $\mathrm{FOL}-\mathrm{BC}-\mathrm{OR}(K B$, goal, $\theta)$ returns a substitution
for each rule in Fetch-Rules-For-Goal( $K B$, goal) do
$(l h s \Rightarrow r h s) \leftarrow$ STANDARDIZE-VARIABLES $(r u l e)$
for each $\theta^{\prime}$ in FOL-BC-AND $(K B$, $l h s$, $\operatorname{UNIFY}(r h s, g o a l, \theta))$ do yield $\theta^{\prime}$
function FOL-BC-AND ( $K B$, goals, $\theta$ ) returns a substitution if $\theta=$ failure then return
else if LENGTH $($ goals $)=0$ then yield $\theta$
else
first, rest $\leftarrow \mathrm{FIRST}($ goals $), \operatorname{REST}($ goals $)$
for each $\theta^{\prime}$ in $\operatorname{FOL}-\mathrm{BC}-\mathrm{Or}(K B, \operatorname{Subst}(\theta$, first $), \theta)$ do
for each $\theta^{\prime \prime}$ in FOL-BC- $\operatorname{AND}\left(K B\right.$, rest, $\left.\theta^{\prime}\right)$ do yield $\theta^{\prime \prime}$

Proof tree constructed by backward chaining t o prove that West is a criminal


## Pseudocode representing the result of compiling the Append predicate

```
procedure APPEND( ax, y, az, continuation)
    trail \leftarrowGLOBAL-TRAIL-POINTER()
    if }ax=[] \mathrm{ and UNIFY(y,az) then CALL(continuation)
    RESET-TRAIL(trail)
    a,x,z\leftarrowNEW-VARIABLE(), NEW-VARIABLE(), NEW-VARIABLE()
    if UNIFY(ax,[a]+x) and UNIFY( }az,[a|z])\mathrm{ then APPEND(x,y,z, continuation)
```


## Finding a path from A to C can lead Prolog into an infinite loop.


(a)

(b)

## Proof that a path exists from A to C.



## A resolution proof that West is a criminal



## A resolution proof that Curiosity killed the cat



Structure of a

## completeness proof for resolution

Any set of sentences $S$ is representable in clausal form


Assume $S$ is unsatisfiable, and in clausal form


Some set $S^{\prime}$ of ground instances is unsatisfiable


There is a resolution proof for the contradiction in $S^{\prime}$

## Knowledge <br> Representation

## The Upper Ontology of the World



## Predicates on time intervals



A schematic view of the object
President (USA) for the early years


## A semantic network

with four objects (John, Mary, 1, and 2) and four categories Relations are denoted by labeled links


## Semantic network

Representation of the logical assertion Fly (Shankar, NewYork, NewDelhi, Yesterday)


## The syntax of descriptions in a subset of the CLASSIC language.

```
    Concept }->\mathrm{ Thing | ConceptName
        And(Concept, ...)
        All(RoleName, Concept)
        AtLeast(Integer, RoleName)
        AtMost(Integer, RoleName)
        Fills(RoleName, IndividualName,...)
        SameAs(Path, Path)
        OneOf(IndividualName,...)
        Path }->\mathrm{ [RoleName,...]
ConceptName }->\mathrm{ Adult | Female| Male|...
    RoleName }->\mathrm{ Spouse | Daughter | Son | ...
```


## Knowledge Graph (KG)

## Knowledge Graph (KG)

- Knowledge Graph (KG)
- A knowledge graph is a multi-relational graph composed of entities and relations, which are regarded as
nodes and different types of edges, respectively (Ji et al., 2021).
- Represents knowledge as concepts (entities) and their relationships (Facts)
- Triple of facts
- SPO: (subject, predicate, object)
- HRT: (head, relation, tail)
- Common Knowledge Graph: DBpedia, YAGO, Wikidata


## Knowledge Graph, Facts, Triple, Embedding

- $G$
- Knowledge graph
- $F$
- Set of facts
- $(h, r, t)$
- Triple of head, relation, and tail
- (h, r, t)
- Embedding of head, relation, and tail


## Knowledge Representation Factual Triple and Knowledge Graph

- Albert Einstein, winner of the 1921 Nobel prize in physics
- The Nobel Prize in Physics 1921 was awarded to Albert Einstein "for his services to Theoretical Physics, and especially for his discovery of the law of the photoelectric effect."


## Knowledge <br> Graph



## Factual Triples in Knowledge Base

$$
(h, r, t)
$$

(Albert Einstein, BornIn, German Empire)
(Albert Einstein, SonOf, Hermann Einstein)
(Albert Einstein, GraduateFrom, University of Zurich)
(Albert Einstein, WinnerOf, Nobel Prize in Physics)
(Albert Einstein, ExpertIn, Physics)
(Nobel Prize in Physics, AwardIn, Physics)
(The theory of relativity, TheoryOf, Physics)
(Albert Einstein, SupervisedBy, Alfred Kleiner)
(Alfred Kleiner, ProfessorOf, University of Zurich)
(The theory of relativity, ProposedBy, Albert Einstein) (Hans Albert Einstein, SonOf, Albert Einstein)

## Entities and Relations in Knowledge Graph

(Albert Einstein, WinnerOf, Nobel Prize in Physics)


## knowledge base and knowledge graph

Factual triples in knowledge base
Entities and relations in knowledge graph
(Albert Einstein, BornIn, German Empire) (Albert Einstein, SonOf, Hermann Einstein) (Albert Einstein, GraduateFrom, University of Zurich) (Albert Einstein, WinnerOf, Nobel Prize in Physics)
(Albert Einstein, ExpertIn, Physics) (Nobel Prize in Physics, AwardIn, Physics) (The theory of relativity, TheoryOf, Physics)
(Albert Einstein, SupervisedBy, Alfred Kleiner) (Alfred Kleiner, ProfessorOf, University of Zurich) (The theory of relativity, ProposedBy, Albert Einstein) (Hans Albert Einstein, SonOf, Albert Einstein)
(Albert Einstein, WinnerOf, Nobel Prize in Physics)


## Categorization of Research on Knowledge Graphs



## Knowledge Graph Completion (KGC) Datasets

| Knowledge Graph <br> Completion (KGC) <br> Dataset | \#Entity | \#Relation | \#Train | \#Valid | \#Test | Reference |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| WN18RR | 40,943 | 11 | 86,835 | 3,034 | 3,134 | Toutanova \& Chen <br> (2015); <br> Zhang et al. (2020) |
| FB15k-237 | 14,541 | 237 | 272,115 | 17,535 | 20,466 | Dettmers et al. (2018); <br> Zhang et al. (2020) |
| YAGO3-10 | 123,182 | 37 | $1,079,040$ | 5,000 | 5,000 | Mahdisoltani et al. <br> (2015); Zhang et al. <br> (2020) |

## Domain-Specific Knowledge Graph

- Domain-Specific Knowledge Graph
- PubMed Knowledge Graph (PKG)
- Extracting biological entities from 29 million PubMed abstracts
- Lynx: Legal Knowledge Graph for Multilingual Compliance Services
- Legal Knowledge Graph (LKG) integrates and links heterogeneous compliance data sources including legislation, case law, standards and other private contracts.


## Lynx: Legal Knowledge Graph for Multilingual Compliance Services



# Automated Planning 

## A PDDL description of an air cargo transportation planning problem

$\operatorname{Init}\left(A t\left(C_{1}, S F O\right) \wedge \operatorname{At}\left(C_{2}, J F K\right) \wedge \operatorname{At}\left(P_{1}, S F O\right) \wedge \operatorname{At}\left(P_{2}, J F K\right)\right.$
$\wedge \operatorname{Cargo}\left(C_{1}\right) \wedge \operatorname{Cargo}\left(C_{2}\right) \wedge \operatorname{Plane}\left(P_{1}\right) \wedge \operatorname{Plane}\left(P_{2}\right)$
$\wedge \operatorname{Airport}(J F K) \wedge \operatorname{Airport}(S F O))$
$\operatorname{Goal}\left(\operatorname{At}\left(C_{1}, J F K\right) \wedge \operatorname{At}\left(C_{2}, S F O\right)\right)$
Action $(\operatorname{Load}(c, p, a)$,
PRECOND: $\operatorname{At}(c, a) \wedge \operatorname{At}(p, a) \wedge \operatorname{Cargo}(c) \wedge \operatorname{Plane}(p) \wedge \operatorname{Airport}(a)$
Effect: $\neg A t(c, a) \wedge \operatorname{In}(c, p))$
Action( Unload $(c, p, a)$,
Precond: $\operatorname{In}(c, p) \wedge \operatorname{At}(p, a) \wedge \operatorname{Cargo}(c) \wedge \operatorname{Plane}(p) \wedge \operatorname{Airport}(a)$
EFFECT: $\operatorname{At}(c, a) \wedge \neg \operatorname{In}(c, p))$
Action(Fly ( $p$, from, to),
Precond: $\operatorname{At}(p$, from $) \wedge \operatorname{Plane}(p) \wedge \operatorname{Airport}($ from $) \wedge \operatorname{Airport}(t o)$
Effect: $\neg A t(p$, from $) \wedge A t(p, t o))$

## The simple spare tire problem

```
\(\operatorname{Init}(\) Tire \((\) Flat \() \wedge \operatorname{Tire}(\) Spare \() \wedge\) At (Flat, Axle \() \wedge\) At(Spare, Trunk \())\)
Goal(At(Spare, Axle))
Action(Remove (obj, loc),
    Precond: At (obj, loc)
    Effect: ᄀ At (obj, loc) \(\wedge\) At (obj, Ground))
Action(PutOn (t, Axle),
    Precond: Tire \((t) \wedge A t(t\), Ground \() \wedge \neg A t(\) Flat, Axle \() \wedge \neg A t(\) Spare, Axle \()\)
    Effect: \(\neg A t(t\), Ground \() \wedge \operatorname{At}(t\), Axle \())\)
Action(LeaveOvernight,
    Precond:
    Effect: \(\neg\) At (Spare, Ground) \(\wedge \neg\) At(Spare, Axle) \(\wedge \neg\) At(Spare, Trunk)
        \(\wedge \neg\) At (Flat, Ground \() \wedge \neg\) At (Flat, Axle \() \wedge \neg\) At (Flat, Trunk) \()\)
```


## Diagram of the blocks-world problem



Start State


Goal State

## A planning problem in the blocks world: building a three-block tower

```
\(\operatorname{Init}(O n(A\), Table \() \wedge O n(B, T a b l e) \wedge O n(C, A)\)
    \(\wedge \operatorname{Block}(A) \wedge \operatorname{Block}(B) \wedge \operatorname{Block}(C) \wedge \operatorname{Clear}(B) \wedge \operatorname{Clear}(C) \wedge \operatorname{Clear}(\) Table \())\)
\(\operatorname{Goal}(O n(A, B) \wedge O n(B, C))\)
Action(Move ( \(b, x, y\) ),
    Precond: On \((b, x) \wedge \operatorname{Clear}(b) \wedge \operatorname{Clear}(y) \wedge \operatorname{Block}(b) \wedge \operatorname{Block}(y) \wedge\)
        \((b \neq x) \wedge(b \neq y) \wedge(x \neq y)\),
    Effect: On \((b, y) \wedge \operatorname{Clear}(x) \wedge \neg \operatorname{On}(b, x) \wedge \neg \operatorname{Clear}(y))\)
Action(MoveToTable (b, \(x\) ),
    Precond: On \((b, x) \wedge \operatorname{Clear}(b) \wedge \operatorname{Block}(b) \wedge \operatorname{Block}(x)\),
    Effect: On \((b\), Table \() \wedge \operatorname{Clear}(x) \wedge \neg O n(b, x))\)
```


## Two approaches to searching for a plan (a)

Forward (progression) search
(b) Backward (regression) search


Two state spaces from planning problems with the ignore-delete-lists heuristic


## Definitions of possible refinements for two high-level actions

```
Refinement(Go(Home, SFO),
    STEPS: [Drive(Home, SFOLongTermParking),
    Shuttle(SFOLongTermParking, SFO)] )
Refinement(Go(Home, SFO),
    Steps: [Taxi(Home, SFO)] )
Refinement(Navigate \(([a, b],[x, y])\),
    PRECOND: \(a=x \wedge b=y\)
    Steps: [] )
Refinement(Navigate \(([a, b],[x, y])\),
    Precond:Connected \(([a, b],[a-1, b])\)
    Steps: \([\) Left, Navigate \(([a-1, b],[x, y])])\)
Refinement(Navigate \(([a, b],[x, y])\),
    Precond:Connected \(([a, b],[a+1, b])\)
    Steps: \([\operatorname{Right}\), Navigate \(([a+1, b],[x, y])])\)
```


## A breadth-first implementation of hierarchical forward planning search

function HIERARCHICAL-SEARCH( problem, hierarchy) returns a solution or failure
frontier $\leftarrow$ a FIFO queue with $[A c t]$ as the only element while true do
if Is-Empty ( frontier) then return failure
plan $\leftarrow \operatorname{POP}($ frontier $) \quad / /$ chooses the shallowest plan in frontier
$h l a \leftarrow$ the first HLA in plan, or null if none
prefix,suffix $\leftarrow$ the action subsequences before and after hla in plan
outcome $\leftarrow \operatorname{RESULT}($ problem.INITIAL, prefix)
if hla is null then / so plan is primitive and outcome is its result if problem.Is-GoAL(outcome) then return plan
else for each sequence in REFINEMENTS(hla, outcome, hierarchy) do add APPEND ( prefix, sequence, suffix) to frontier

## Schematic examples of reachable sets



Goal achievement for high-level plans with approximate descriptions


## A hierarchical planning algorithm

```
function ANGELIC-SEARCH(problem, hierarchy, initialPlan) returns solution or fail
    frontier }\leftarrow\mathrm{ a FIFO queue with initialPlan as the only element
    while true do
        if Empty?(frontier) then return fail
        plan \leftarrowPOP(frontier) // chooses the shallowest node in frontier
        if REACH }\mp@subsup{}{}{+}\mathrm{ (problem.Initial, plan) intersects problem.GoAL then
            if plan is primitive then return plan // REACH
            guaranteed }\leftarrow\mp@subsup{\textrm{REACH}}{}{-}(\mathrm{ problem.InitiAL, plan ) }\cap\mathrm{ problem.GoAL
            if guaranteed }\not={}\mathrm{ and MAKING-Progress(plan, initialPlan) then
            finalState }\leftarrow\mathrm{ any element of guaranteed
            return DECOMPOSE(hierarchy, problem.InITIAL, plan, finalState)
            hla}\leftarrow\mathrm{ some HLA in plan
            prefix,suffix \leftarrowt the action subsequences before and after hla in plan
            outcome \leftarrow & RESULT(problem.InITIAL, prefix)
            for each sequence in RefinEmENTS(hla,outcome, hierarchy) do
            frontier }\leftarrow\mathrm{ Insert(APPEND(prefix, sequence, suffix), frontier)
```


## A hierarchical planning algorithm Decompose solution

function DECOMPOSE (hierarchy, $s_{0}$, plan, $s_{f}$ ) returns a solution

```
solution }\leftarrow\mathrm{ an empty plan
while plan is not empty do
    action }\leftarrow\mathrm{ REMOVE-LAST(plan)
    si}\leftarrow\mp@code{a state in REACH}\mp@subsup{}{}{-}(\mp@subsup{s}{0}{},\mathrm{ plan ) such that }\mp@subsup{s}{f}{}\in\mp@subsup{\operatorname{REACH}}{}{-}(\mp@subsup{s}{i}{},\mathrm{ action }
    problem}\leftarrow\mathrm{ a problem with INITIAL = si}\mathrm{ and GOAL = sf
    solution }\leftarrow\mathrm{ APPEND(ANGELIC-SEARCH(problem, hierarchy,action), solution)
    sf}\leftarrow\mp@subsup{s}{i}{
return solution
```


## At first, the sequence "whole plan" is expected to get the agent from $S$ to $G$



## A job-shop scheduling problem for assembling two cars, with resource constraints

$\begin{aligned} & \text { Jobs }(\{\text { AddEngine } 1 \prec \text { AddWheels } 1 \\ &\text { Ad Inspect1 }\}, \\ & \text { AdEngine } 2 \prec \text { AddWheels2 } \\ &\text { Inspect } 2\})\end{aligned}$

Resources(EngineHoists(1), WheelStations(1), Inspectors(e2), LugNuts(500))
Action(AddEngine1, DURATION:30,
USE:EngineHoists(1))
Action(AddEngine2, Duration:60,
USE:EngineHoists(1))
Action(AddWheels1, Duration:30,
Consume:LugNuts(20), Use:WheelStations(1))
Action(AddWheels2, Duration:15,
Consume:LugNuts(20), Use: WheelStations(1))
Action( Inspect $_{i}$, DURATION:10,
UsE:Inspectors(1))

## A representation of the temporal constraints for the job-shop scheduling problem



Source: Stuart Russell and Peter Norvig (2020), Artificial Intelligence: A Modern Approach, 4th Edition, Pearson

## A solution to the

## job-shop scheduling problem



## Artificial Intelligence: Uncertain Knowledge and Reasoning

## Artificial Intelligence:

4. Uncertain Knowledge and Reasoning

- Quantifying Uncertainty
- Probabilistic Reasoning
- Probabilistic Reasoning over Time
- Probabilistic Programming
- Making Simple Decisions
- Making Complex Decisions
- Multiagent Decision Making


## Quantifying Uncertainty

# DT-Agent <br> A Decision-Theoretic Agent that Selects Rational Actions 

function DT-AGENT (percept) returns an action persistent: belief_state, probabilistic beliefs about the current state of the world action, the agent's action
update belief_state based on action and percept calculate outcome probabilities for actions, given action descriptions and current belief_state select action with highest expected utility
given probabilities of outcomes and utility information return action

## Agent 1 has inconsistent beliefs

## Proposition Agent 1's Agent2 Agent1 Agent 1 payoffs for each outcome <br> belief bets bets $a, b a, \neg b \neg a, b \neg a, \neg b$

$$
\begin{array}{cccccccc}
a & 0.4 & \$ 4 o n a & \$ 6 \text { on } \sim a & -\$ 6 & -\$ 6 & \$ 4 & \$ 4 \\
b & 0.3 & \$ 3 \text { on } b & \$ 70 n-b & -\$ 7 & \$ 3 & -\$ 7 & \$ 3 \\
a \vee b & 0.8 & \$ 2 \text { on } 7(a \vee b) & \$ 8 \text { on } a \vee b & \$ 2 & \$ 2 & \$ 2 & -\$ 8 \\
& & & & -\$ 11 & -\$ 1 & -\$ 1 & -\$ 1
\end{array}
$$

## A full joint distribution for the Toothache, Cavity, Catch world

|  | toothache |  | -toothache |  |
| :---: | :---: | :---: | :---: | :---: |
|  | catch | ratch | catch | רcatch |
| cavity | 0.108 | 0.012 | 0.072 | 0.008 |
| racuity | 0.016 | 0.064 | 0.144 | 0.576 |

## Weather and Dental problems are independent



## Coin flips are independent



## Probabilistic Reasoning

## A Simple Bayesian Network

Weather is independent to the other three variables. Toothache and Catch are conditionally independent, given Cavity.


## A Typical Bayesian Network

Topology and the Conditional Probability Tables (CPTs)


## Conditional Probability Table for P(Fever |Cold, Flu, Malaria)

| Cold | Flu | Malaria $P($ fever $\mid \cdot)$ | $P(\neg$ fever $\mid \cdot)$ |  |
| :---: | :---: | :---: | :--- | :--- |
| $f$ | $f$ | $f$ | 0.0 | 1.0 |
| $f$ | $f$ | $t$ | 0.9 | $\mathbf{0 . 1}$ |
| $f$ | $t$ | $f$ | 0.8 | $\mathbf{0 . 2}$ |
| $f$ | $t$ | $t$ | 0.98 | $0.02=0.2 \times 0.1$ |
| $t$ | $f$ | $f$ | 0.4 | $\mathbf{0 . 6}$ |
| $t$ | $f$ | $t$ | 0.94 | $0.06=0.6 \times 0.1$ |
| $t$ | $t$ | $f$ | 0.88 | $0.12=0.6 \times 0.2$ |
| $t$ | $t$ | $t$ | 0.988 | $0.012=0.6 \times 0.2 \times 0.1$ |

## A Simple Network

## with discrete variables (Subsidy and Buys) and continuous variables (Harvest and Cost )



## Probability distribution

## over Cost as a function of Harvest size



## A normal (Gaussian) distribution for the cost threshold


(a)

(b)

Expit and Probit models for the probability of buys given cost

## A Bayesian Network

for evaluating car insurance applications


## The structure of the expression



## The Enumeration Algorithm for Exact Inference in Bayes Nets

function EnUMERATION-ASK $(X, \mathbf{e}, b n)$ returns a distribution over $X$ inputs: $X$, the query variable
$\mathbf{e}$, observed values for variables $\mathbf{E}$ $b n$, a Bayes net with variables vars
$\mathbf{Q}(X) \leftarrow$ a distribution over $X$, initially empty
for each value $x_{i}$ of $X$ do
$\mathbf{Q}\left(x_{i}\right) \leftarrow$ Enumerate-AlL $\left(\right.$ vars, $\left.\mathbf{e}_{x_{i}}\right)$
where $\mathbf{e}_{x_{i}}$ is $\mathbf{e}$ extended with $X=x_{i}$
return Normalize $(\mathbf{Q}(X))$
function EnUMERATE-ALL(vars, e) returns a real number
if Empty? (vars) then return 1.0
$V \leftarrow \mathrm{FIRST}($ vars $)$
if $V$ is an evidence variable with value $v$ in $\mathbf{e}$
then return $P(v \mid$ parents $(V)) \times$ Enumerate-All(Rest(vars), e)
else return $\sum_{v} P(v \mid \operatorname{parents}(V)) \times$ Enumerate-All(Rest(vars), $\left.\mathbf{e}_{v}\right)$
where $\mathbf{e}_{v}$ is $\mathbf{e}$ extended with $V=v$

## Pointwise Multiplication $\mathrm{f}(X, Y) \times \mathrm{g}(Y, Z)=\mathrm{h}(X, Y, Z)$

| $X$ | $Y$ | $\mathbf{f}(X, Y)$ | $Y$ | $Z$ | $\mathbf{g}(Y, Z)$ | $X$ | $Y$ | $Z$ | $\mathbf{h}(X, Y, Z)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $t$ | $t$ | .3 | $t$ | $t$ | .2 | $t$ | $t$ | $t$ | $.3 \times .2=.06$ |
| $t$ | $f$ | .7 | $t$ | $f$ | .8 | $t$ | $t$ | $f$ | $.3 \times .8=.24$ |
| $f$ | $t$ | .9 | $f$ | $t$ | .6 | $t$ | $f$ | $t$ | $.7 \times .6=.42$ |
| $f$ | $f$ | .1 | $f$ | $f$ | .4 | $t$ | $f$ | $f$ | $.7 \times .4=.28$ |
|  |  |  |  |  |  | $f$ | $t$ | $t$ | $.9 \times .2=.18$ |
|  |  |  |  |  |  | $f$ | $t$ | $f$ | $.9 \times .8=.72$ |
|  |  |  |  |  |  | $f$ | $f$ | $t$ | $.1 \times .6=.06$ |
|  |  |  |  |  |  | $f$ | $f$ | $.1 \times .4=.04$ |  |

## The Variable Elimination Algorithm for Exact Inference in Bayes Nets

function ELIMINATION- $\operatorname{AsK}(X, \mathbf{e}, b n)$ returns a distribution over $X$ inputs: $X$, the query variable
$\mathbf{e}$, observed values for variables $\mathbf{E}$
$b n$, a Bayesian network with variables vars
factors $\leftarrow[]$
for each $V$ in ORDER(vars) do
factors $\leftarrow[\operatorname{MAKE}-\mathrm{FACTOR}(V, \mathbf{e})]+$ factors
if $V$ is a hidden variable then factors $\leftarrow$ Sum-Out ( $V$, factors)
return Normalize(Pointwise-Product(factors))

# Bayes Net Encoding 

of the 3-CNF (Conjunctive Normal Form) Sentence $(W \vee X V Y) \wedge(\neg W \vee Y V Z) \wedge(X \vee Y \vee \neg Z)$


## Multiply Connected Network

(b) A clustered equivalent


## A Sampling Algorithm that generates events from a Bayesian network

function PRIIoR-SAMPLE( $b n$ ) returns an event sampled from the prior specified by $b n$ inputs: bn, a Bayesian network specifying joint distribution $\mathbf{P}\left(X_{1}, \ldots, X_{n}\right)$
$\mathbf{x} \leftarrow$ an event with $n$ elements
for each variable $X_{i}$ in $X_{1}, \ldots, X_{n}$ do
$\mathbf{x}[i] \leftarrow \operatorname{arandom}$ sample from $\mathbf{P}\left(X_{i} \mid \operatorname{parents}\left(X_{i}\right)\right)$
return x

## The Rejection-Sampling Algorithm

## for answering queries given evidence in a Bayesian network

function REJECTION-SAMPLING $(X, \mathbf{e}, b n, N)$ returns an estimate of $\mathbf{P}(X \mid \mathbf{e})$ inputs: $X$, the query variable
$\mathbf{e}$, observed values for variables $\mathbf{E}$
$b n$, a Bayesian network
$N$, the total number of samples to be generated
local variables: $\mathbf{C}$, a vector of counts for each value of $X$, initially zero
for $j=1$ to $N$ do
$\mathbf{x} \leftarrow \operatorname{PRIOR}-\operatorname{SAMPLE}(b n)$
if $\mathbf{x}$ is consistent with $\mathbf{e}$ then
$\mathbf{C}[j] \leftarrow \mathbf{C}[j]+1$ where $x_{j}$ is the value of $X$ in $\mathbf{x}$
return Normalize( $\mathbf{C}$ )

## The Likelihood-Weighting Algorithm for inference in Bayesian networks

function Likelihood-Weighting $(X, \mathbf{e}, b n, N)$ returns an estimate of $\mathbf{P}(X \mid \mathbf{e})$ inputs: $X$, the query variable
$\mathbf{e}$, observed values for variables $\mathbf{E}$
$b n$, a Bayesian network specifying joint distribution $\mathbf{P}\left(X_{1}, \ldots, X_{n}\right)$
$N$, the total number of samples to be generated
local variables: $\mathbf{W}$, a vector of weighted counts for each value of $X$, initially zero

$$
\text { for } j=1 \text { to } N \text { do }
$$

$\mathbf{x}, w \leftarrow$ WEIGHTED-SAMPLE $(b n, \mathbf{e})$
$\mathbf{W}[j] \leftarrow \mathbf{W}[j]+w$ where $x_{j}$ is the value of $X$ in $\mathbf{x}$ return $\operatorname{Normalize}(\mathbf{W})$
function Weighted-SAMPLE $(b n, \mathbf{e})$ returns an event and a weight
$w \leftarrow 1 ; \mathbf{x} \leftarrow$ an event with $n$ elements, with values fixed from $\mathbf{e}$
for $i=1$ to $n$ do
if $X_{i}$ is an evidence variable with value $x_{i j}$ in $\mathbf{e}$
then $w \leftarrow w \times P\left(X_{i}=x_{i j} \mid\right.$ parents $\left.\left(X_{i}\right)\right)$
else $\mathbf{x}[i] \leftarrow$ a random sample from $\mathbf{P}\left(X_{i} \mid\right.$ parents $\left.\left(X_{i}\right)\right)$
return $\mathbf{x}, w$

## Performance of rejection sampling and likelihood weighting on the insurance network



## The Gibbs Sampling Algorithm for approximate inference in Bayes nets

function $\operatorname{Gibbs-Ask}(X, \mathbf{e}, b n, N)$ returns an estimate of $\mathbf{P}(X \mid \mathbf{e})$
local variables: $\mathbf{C}$, a vector of counts for each value of $X$, initially zero
$\mathbf{Z}$, the nonevidence variables in $b n$
$\mathbf{x}$, the current state of the network, initialized from $\mathbf{e}$
initialize $\mathbf{x}$ with random values for the variables in $\mathbf{Z}$ for $k=1$ to $N$ do
choose any variable $Z_{i}$ from $\mathbf{Z}$ according to any distribution $\rho(i)$
set the value of $Z_{i}$ in $\mathbf{x}$ by sampling from $\mathbf{P}\left(Z_{i} \mid m b\left(Z_{i}\right)\right)$
$\mathbf{C}[j] \leftarrow \mathbf{C}[j]+1$ where $x_{j}$ is the value of $X$ in $\mathbf{x}$
return Normalize( $\mathbf{C}$ )

## The States and Transition Probabilities of the Markov Chain

for the query $\mathbf{P}($ Rain I Sprinkler $=$ true, WetGrass $=$ true $)$

(a)

(b)

Transition Probabilities
when the CPT for Rain constrains it to have the same value as Cloudy

## Performance of Gibbs sampling compared to likelihood weighting on the car insurance network


(a)
for the standard query on PropertyCost

(b)
for the case where the output variables are observed and Age is the query variable

## A Causal Bayesian Network

## representing cause-effect relations among five variables



The network after performing the action "turn Sprinkler on."

## Probabilistic

 Reasoning over Time
## Bayesian network structure

## corresponding to a First-order Markov Process

 with state defined by the variables $X t$.

## Bayesian Network Structure and Conditional Distributions describing the umbrella world



## Smoothing computes $P\left(X_{k} \mid e_{1: t}\right)$

 the posterior distribution of the state at some past time $k$ given a complete sequence of observations from 1 to t .

## The Forward-Backward Algorithm for Smoothing

function FORWARD-BACKWARD(ev, prior) returns a vector of probability distributions
inputs: ev, a vector of evidence values for steps $1, \ldots, t$
prior, the prior distribution on the initial state, $\mathbf{P}\left(\mathbf{X}_{0}\right)$
local variables: $\mathbf{f v}$, a vector of forward messages for steps $0, \ldots, t$
b, a representation of the backward message, initially all 1 s
$\mathbf{s v}$, a vector of smoothed estimates for steps $1, \ldots, t$
$\mathbf{f v}[0] \leftarrow$ prior
for $i=1$ to $t$ do
$\mathbf{f v}[i] \leftarrow \operatorname{FORWARD}(\mathbf{f v}[i-1], \mathbf{e v}[i])$
for $i=t$ down to 1 do
$\mathbf{s v}[i] \leftarrow \operatorname{NORMALIZE}(\mathbf{f v}[i] \times \mathbf{b})$
$\mathbf{b} \leftarrow \operatorname{BACKWARD}(\mathbf{b}, \mathbf{e v}[i])$
return sv

## Possible state sequences for Rain $_{t}$ can

be viewed as paths through a graph of the possible states
at each time step
(a)


Umbrella ${ }_{t}$
true true
false
true
true

$\mathbf{m}_{1: 0}$
$\mathbf{m}_{1: 1}$
$\mathbf{m}_{1: 2}$
$\mathbf{m}_{1: 3}$
$\mathbf{m}_{1: 4}$
$\mathbf{m}_{1: 5}$
Operation of the Viterbi algorithm for the umbrella observation sequence [true, true, false, true, true]

## Algorithm for Smoothing with a Fixed Time Lag of d Step

function Fixed-Lag-Smoothing $\left(e_{t}, h m m, d\right)$ returns a distribution over $\mathbf{X}_{t-d}$
inputs: $e_{t}$, the current evidence for time step $t$
$h m m$, a hidden Markov model with $S \times S$ transition matrix $\mathbf{T}$ $d$, the length of the lag for smoothing
persistent: $t$, the current time, initially 1
$\mathbf{f}$, the forward message $\mathbf{P}\left(X_{t} \mid e_{1: t}\right)$, initially $h m m$.Prior
B, the $d$-step backward transformation matrix, initially the identity matrix
$e_{t-d: t}$, double-ended list of evidence from $t-d$ to $t$, initially empty
local variables: $\mathbf{O}_{t-d}, \mathbf{O}_{t}$, diagonal matrices containing the sensor model information
add $e_{t}$ to the end of $e_{t-d: t}$
$\mathbf{O}_{t} \leftarrow$ diagonal matrix containing $\mathbf{P}\left(e_{t} \mid X_{t}\right)$
if $t>d$ then
$\mathbf{f} \leftarrow \operatorname{FORWARD}\left(\mathbf{f}, e_{t-d}\right)$
remove $e_{t-d-1}$ from the beginning of $e_{t-d: t}$
$\mathbf{O}_{t-d} \leftarrow$ diagonal matrix containing $\mathbf{P}\left(e_{t-d} \mid X_{t-d}\right)$
$\mathbf{B} \leftarrow \mathbf{O}_{t-d}^{-1} \mathbf{T}^{-1} \mathbf{B T O}_{t}$
else $\mathbf{B} \leftarrow \mathbf{B T O}_{t}$
$t \leftarrow t+1$
if $t>d+1$ then return $\operatorname{Normalize}(\mathbf{f} \times$ B1) else return null

## Specification of the prior, transition model, and sensor model for the umbrella DBN



## A DBN fragment

## the sensor status variable required for modeling persistent failure of the battery sensor

| $B_{0}$ | $P\left(B_{1}\right)$ |
| :---: | :---: |
| $t$ | 1.000 |
| $f$ | 0.001 |


(a)

## Unrolling a

## Dynamic Bayesian Network



## The Particle Filtering Algorithm

function Particle-Filtering $(\mathbf{e}, N, d b n)$ returns a set of samples for the next time step inputs: e, the new incoming evidence
$N$, the number of samples to be maintained
$d b n$, a DBN defined by $\mathbf{P}\left(\mathbf{X}_{0}\right), \mathbf{P}\left(\mathbf{X}_{1} \mid \mathbf{X}_{0}\right)$, and $\mathbf{P}\left(\mathbf{E}_{1} \mid \mathbf{X}_{1}\right)$
persistent: $S$, a vector of samples of size $N$, initially generated from $\mathbf{P}\left(\mathbf{X}_{0}\right)$ local variables: $W$, a vector of weights of size $N$

$$
\begin{array}{ll}
\text { for } i=1 \text { to } N \text { do } & \\
& S[i] \leftarrow \operatorname{sample} \text { from } \mathbf{P}\left(\mathbf{X}_{1} \mid \mathbf{X}_{0}=S[i]\right) \\
& \text { // step 1 } \\
W[i] \leftarrow \mathbf{P}\left(\mathbf{e} \mid \mathbf{X}_{1}=S[i]\right) & \text { //step } 2
\end{array}
$$

$S \leftarrow$ Weighted-Sample-With-Replacement $(N, S, W) \quad / /$ step 3 return $S$

## The Particle Filtering Update Cycle for the Umbrella DBN


(a) Propagate

(b) Weight

Rain $_{t+1}$

(c) Resample

## A Dynamic Bayes Net

for simultaneous localization and mapping in the stochastic-dirt vacuum world


## Probabilistic Programming

## Possible Worlds

for a language with two constant symbols, R and J


## Bayes Net for a Single customer C1

recommending a single book B1. Honest(C1) is Boolean


Bayes net with two customers and two books

## Bayes Net

for the book recommendation when Author(B2) is unknown


## One particular world for the book recommendation OUPM

| Variable | Value | Probability |
| :---: | :---: | :---: |
| \#Customer | 2 | 0.3333 |
| \#Book | 3 | 0.3333 |
| Honest $_{\langle\text {Customer, ,1] }}$ | true | 0.99 |
| Honest ${ }_{\text {(Customer, ,2) }}$ | false | 0.01 |
| Kindness $\langle$ Customer, ,1> | 4 | 0.3 |
| Kindness CCustomer, ,2) $^{\text {, }}$ | 1 | 0.1 |
| Quality ${ }_{\text {(Book, ,1> }}$ | 1 | 0.05 |
| Quality ${ }_{\text {(Book, ,2> }}$ | 3 | 0.4 |
| Quality ${ }_{\text {(Book, ,3> }}$ | 5 | 0.15 |
| \#LoginID $\langle$ Owner, $\langle$ Customer, ,1>> | 1 | 1.0 |
| \#LoginID $\langle$ Owner, <Customer, ,2〉\} | 2 | 0.25 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,1>>,1\}, } \text { Book, }, 1\rangle^{\text {, }}$ | 2 | 0.5 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,1>>,1), } \text { Book, }, 2\rangle^{\text {, }}$ | 4 | 0.5 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,1>>,1\}, } \text { Book, }, 3\rangle^{\text {R }}$ | 5 | 0.5 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,2\}>,1\}, } \text { Book, }, 1\rangle^{\text {, }}$ | 5 | 0.4 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,2\}>,1\}, } \text { Book, }, 2\rangle^{\text {, }}$ | 5 | 0.4 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,2\}>,1\}, } \text { Book, }, 3\rangle^{\text {, }}$ | 1 | 0.4 |
| Recommendation ${ }_{\text {LooginID, },\langle\text { Owner, }\langle\text { Customer, ,2\}}\rangle, 2\rangle,\langle\text { Book, }, 1\rangle}$ | 5 | 0.4 |
|  | 5 | 0.4 |
| Recommendation $\left\langle\right.$ LoginID, $\left\langle\right.$ Owner, $\langle\text { Customer, ,2ो>,2), } \text { Book, }, 3\rangle^{\text {a }}$ | 1 | 0.4 |

## An OUPM for

## Citation Information Extraction

```
type Researcher, Paper, Citation
random String Name(Researcher)
random String Title(Paper)
random Paper PubCited(Citation)
random String Text(Citation)
random Boolean Professor(Researcher)
origin Researcher Author(Paper)
#Researcher ~ OM(3,1)
Name(r) ~ NamePrior()
Professor(r) ~ Boolean(0.2)
#Paper (Author =r) ~ if Professor (r) then OM (1.5,0.5) else OM (1,0.5)
Title(p) ~ PaperTitlePrior()
CitedPaper(c) ~ UniformChoice({Paper p})
Text(c) ~ HMMGrammar(Name(Author(CitedPaper(c))),Title(CitedPaper(c)))
```


# Making Simple <br> Decisions 

## Nontransitive preferences $A>B>C>A$

can result in irrational behavior:
a cycle of exchanges each costing one cent

(a)

is equivalent to

(b)

The decomposability axiom

## The Utility of Money


(a)

(b)

## Unjustified optimism

## caused by choosing the best of $k$ options



## Strict dominance

## (a) Deterministic (b) Uncertain


(a)

(b)

## Stochastic dominance



## A decision network for the airport-siting problem



## A simplified representation of the airport-siting problem



## Making <br> Complex Decisions

## A dynamic decision network

for a mobile robot with state variables for battery level, charging status, location, and velocity, and action variables for the left and right wheel motors and for charging.


## The game of Tetris The DDN for the Tetris MDP


(a)

(b)

## The Value Iteration Algorithm for calculating utilities of states

function VALUE-ITERATION $(m d p, \epsilon)$ returns a utility function inputs: $m d p$, an MDP with states $S$, actions $A(s)$, transition model $P\left(s^{\prime} \mid s, a\right)$, rewards $R\left(s, a, s^{\prime}\right)$, discount $\gamma$
$\epsilon$, the maximum error allowed in the utility of any state
local variables: $U, U^{\prime}$, vectors of utilities for states in $S$, initially zero $\delta$, the maximum relative change in the utility of any state
repeat
$U \leftarrow U^{\prime} ; \delta \leftarrow 0$
for each state $s$ in $S$ do
$U^{\prime}[s] \leftarrow \max _{a \in A(s)} \operatorname{Q-VALUE}(m d p, s, a, U)$
if $\left|U^{\prime}[s]-U[s]\right|>\delta$ then $\delta \leftarrow\left|U^{\prime}[s]-U[s]\right|$
until $\delta \leq \epsilon(1-\gamma) / \gamma$
return $U$

Stuart Russell and Peter Norvig (2020),

## Artificial Intelligence: A Modern Approach,

4th Edition, Pearson


## Artificial Intelligence: A Modern Approach (AIMA)

- Artificial Intelligence: A Modern Approach (AIMA)
- http://aima.cs.berkeley.edu/
- AIMA Python
- http://aima.cs.berkeley.edu/python/readme.html
- https://github.com/aimacode/aima-python
- Logic, KB Agent
- http://aima.cs.berkeley.edu/python/logic.html
- Probability Models (DTAgent)
- http://aima.cs.berkeley.edu/python/probability.html
- Markov Decision Processes (MDP)
- http://aima.cs.berkeley.edu/python/mdp.html


## Artificial Intelligence: A Modern Approach (AIMA)



## Artificial Intelligence: A Modern Approach, 4th US ed.

## by Stuart Russell and Peter Norvig

The authoritative, most-used AI textbook, adopted by over 1500 schools.
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## AIMA Code


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| Norvig's "Artificial Intelligence - A Modern Approach" |  |
| § 740 | \& 386 |


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| :---: | :---: |
| Javascript visualization of algorithms from Russell And Norvig's "Artificial Intelligence - A Modern Approach" | Common Lisp implementation of algorithms from Russell And Norvig's "Artificial Intelligence - A Modern Approach" |
| JavaScript $\begin{gathered} \\ 495 \\ \text { ¢\% }\end{gathered}$ | Common Lisp $\widehat{342}$ ย๊95 |

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- Java § 1.4k ย゚ 767
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- HTML 611 ย゚353


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Natural Language Processing


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https://colab.research.google.com/drive/1FEG6DnGvwfUbeo4zJ1zTunjMaf2RkCrT

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## Summary

- Knowledge and Reasoning
- Logical Agents
- First-Order Logic
- Inference in First-Order Logic
- Knowledge Representation
- Knowledge Graph (KG)
- Uncertain Knowledge and Reasoning
- Quantifying Uncertainty
- Probabilistic Reasoning
- Making Complex Decisions


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